

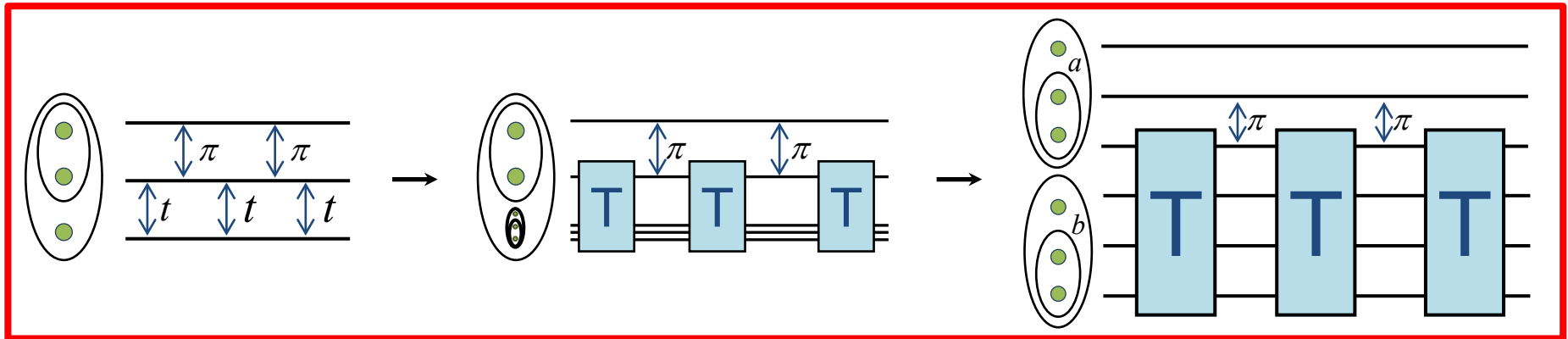


# Designing Two-Qubit Gates for Exchange-Only Quantum Computation



Nick Bonesteel

Dept. of Physics & NHMFL,  
Florida State University



Work done in collaboration with:

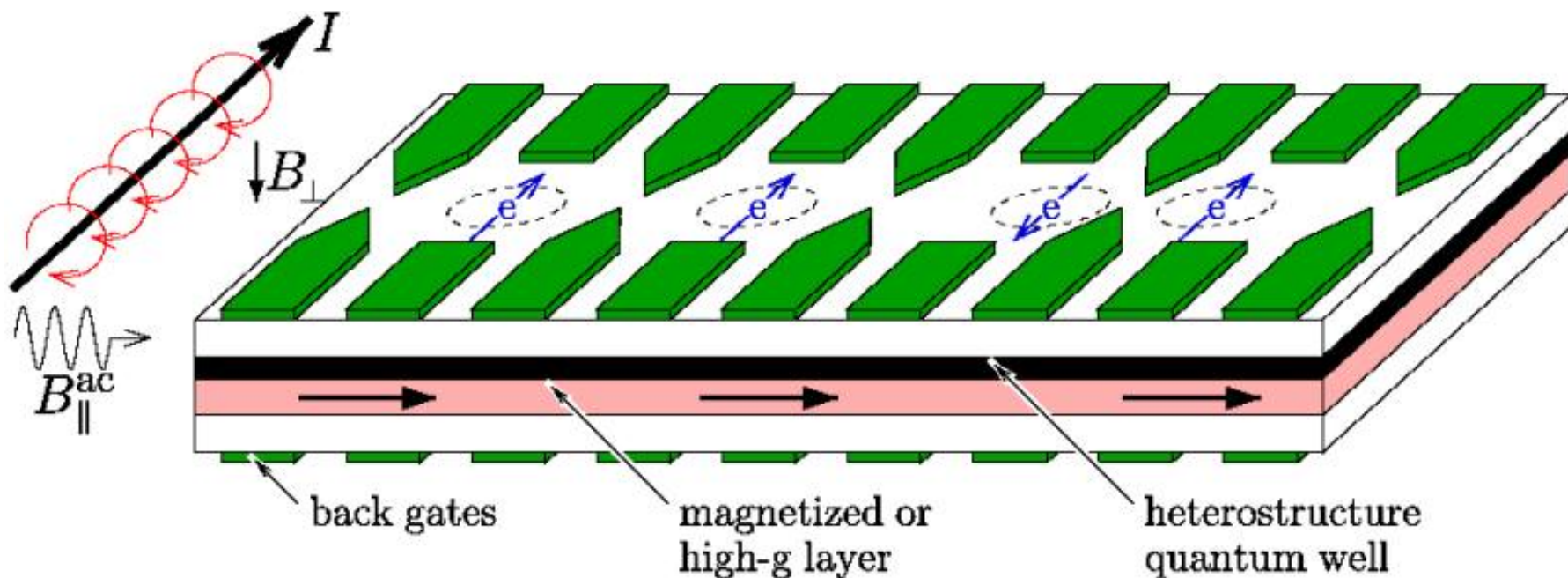
Daniel Zeuch, FSU → Peter Gruenberg Institut, Research Center Juelich

Zeuch, NEB, Phys. Rev. A **98**, 010303 (2016)

Zeuch, NEB, Phys. Rev. B **102**, 075311 (2020)

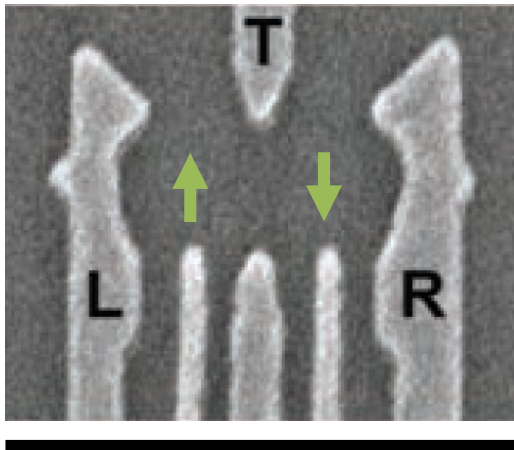
# Early Vision of a Solid State Quantum Computer

Loss & DiVincenzo, Phys. Rev. B (1998)



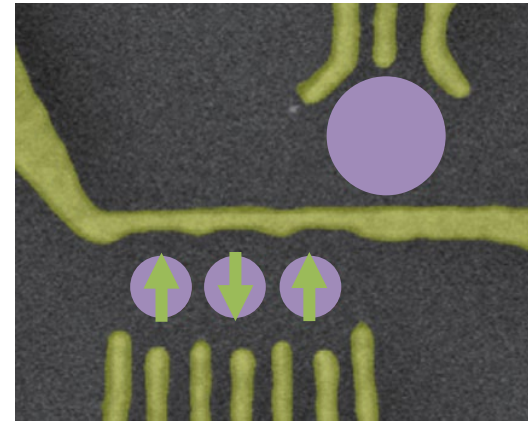
# Decades of Slow Steady Progress

Petta *et al.*, *Science* (2005)



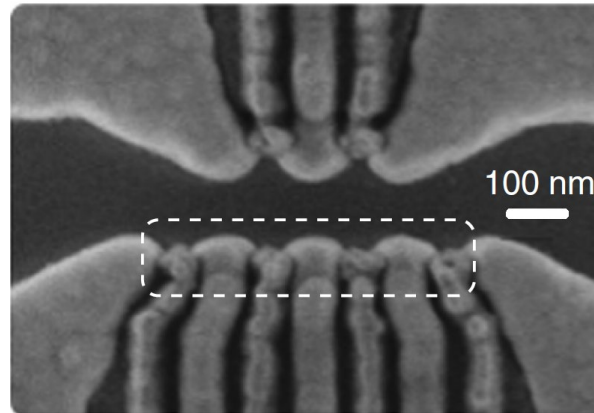
1  $\mu\text{m}$

Medford *et al.*, *Nature Nanotechnology* (2013)



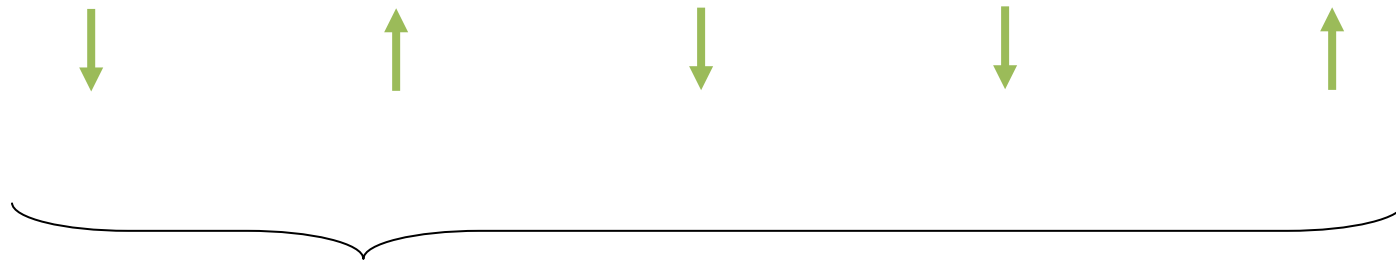
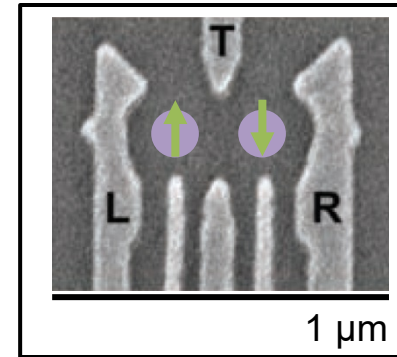
1  $\mu\text{m}$

Andrews *et al.*, *Nature Nanotechnology* (2019)



# Basic Idea

- Use electron spins as qubits

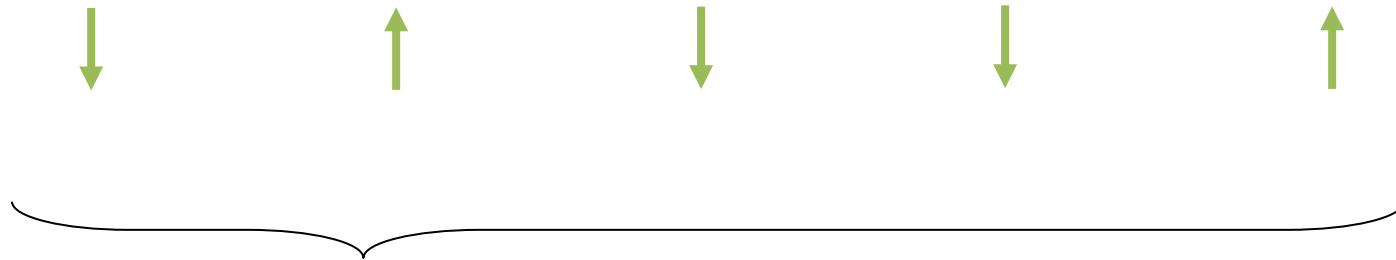
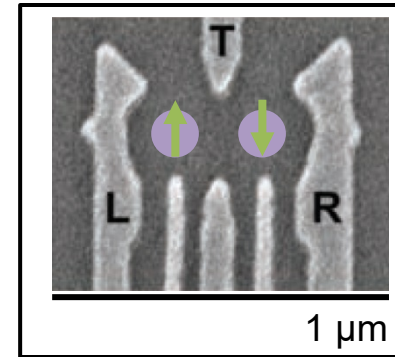


spin-1/2 chain: electrons  
in quantum dots

# Exchange-Based QC

- Quantum gates through spin exchange

$$H_i = J \mathbf{S}_i \cdot \mathbf{S}_{i+1}$$

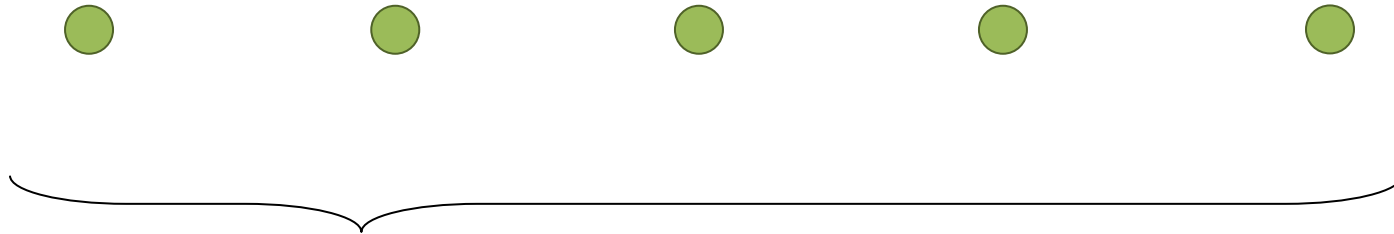
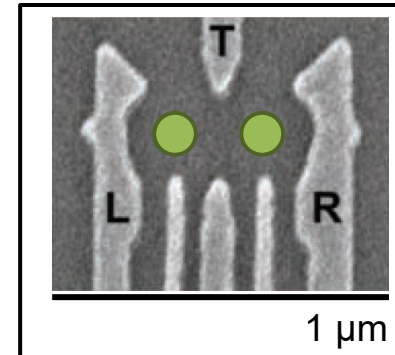


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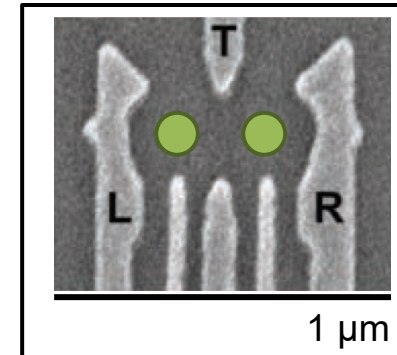
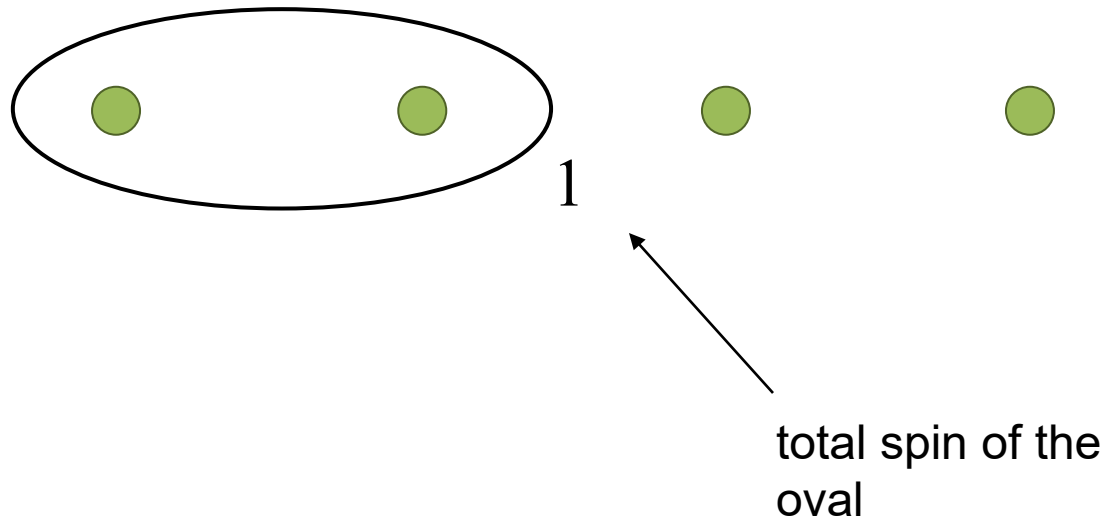


spin-1/2 chain: electrons  
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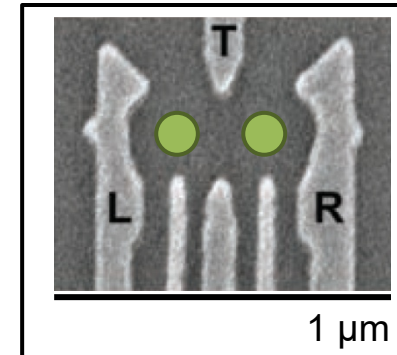
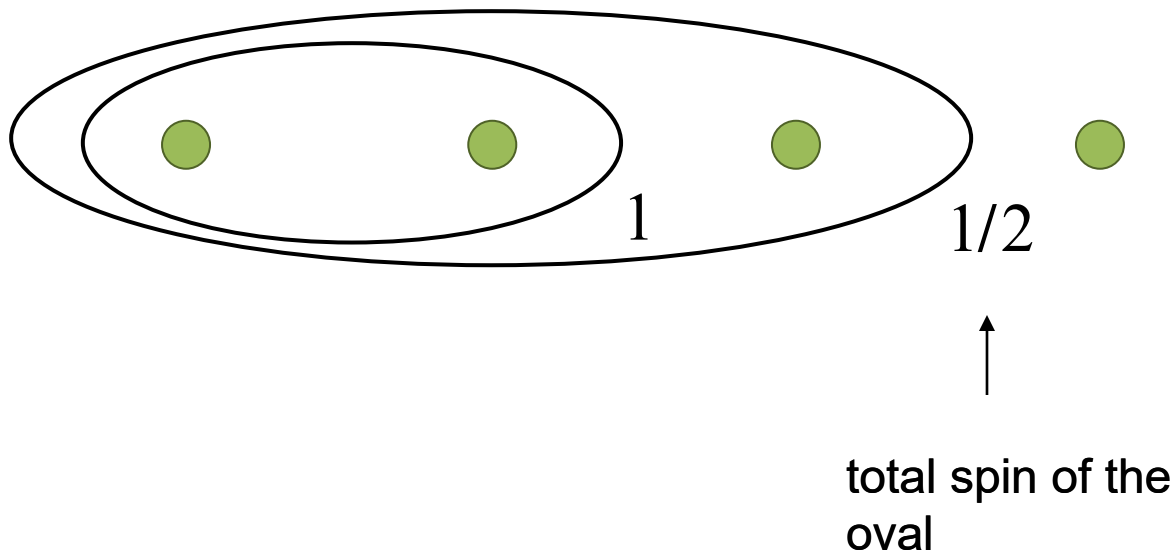
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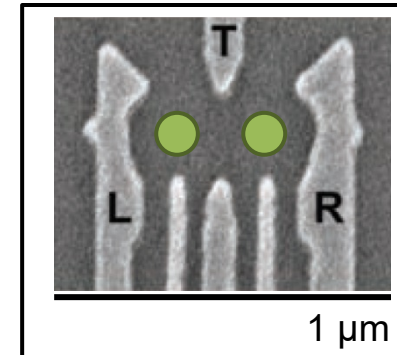
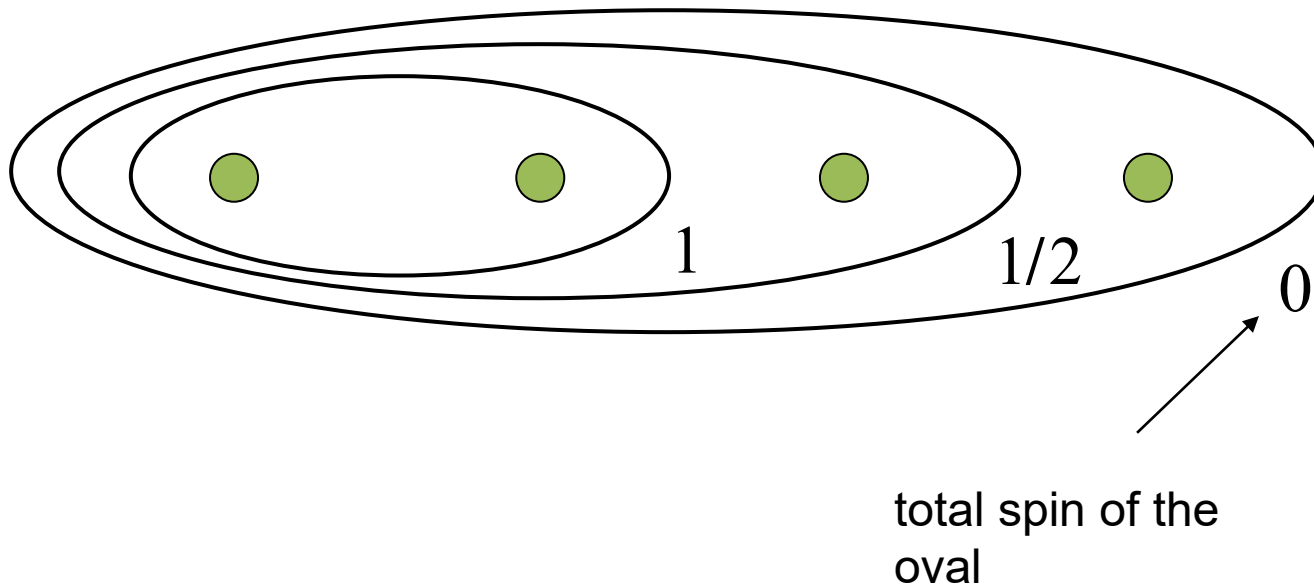




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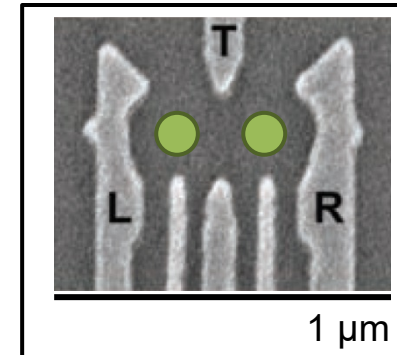
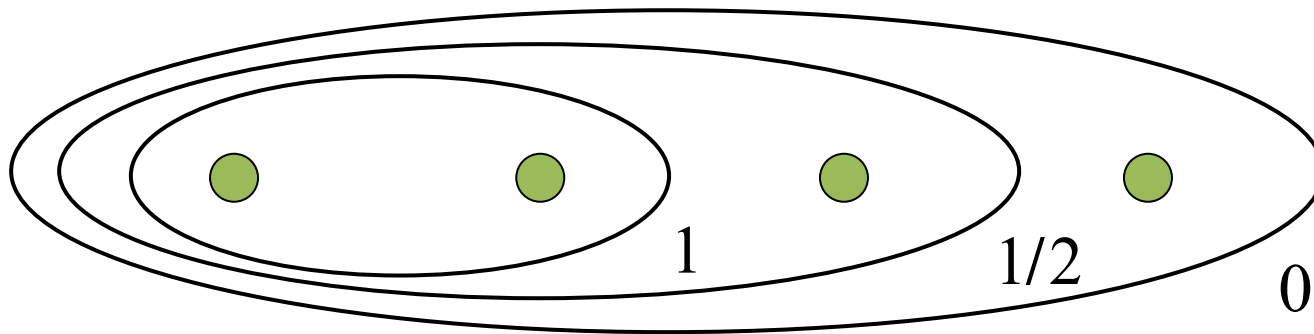
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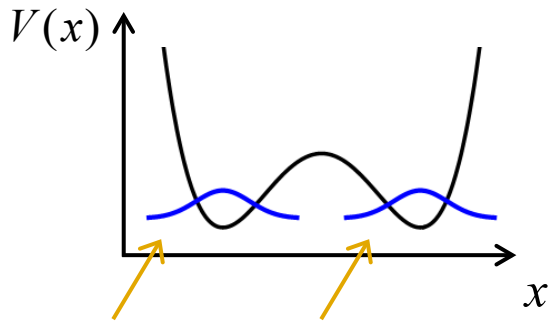
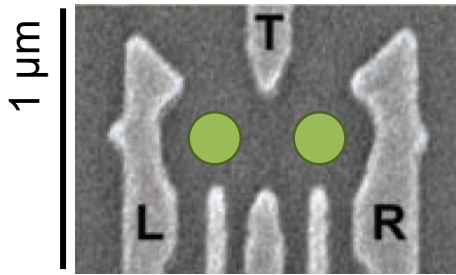
$$H_i = J \mathbf{S}_i \cdot \mathbf{S}_{i+1}$$



$$s_1 \otimes s_2 = |s_1 - s_2|, |s_1 - s_2 + 1|, \dots, s_1 + s_2$$

# Controlling Exchange

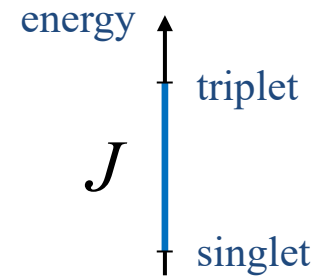
Petta *et al.*, *Science* (2005)



Electron wave functions in quantum dot potential  $V(x)$

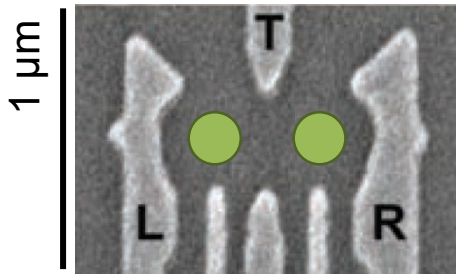
- Exchange Hamiltonian

$$H = J \mathbf{S}_1 \cdot \mathbf{S}_2$$



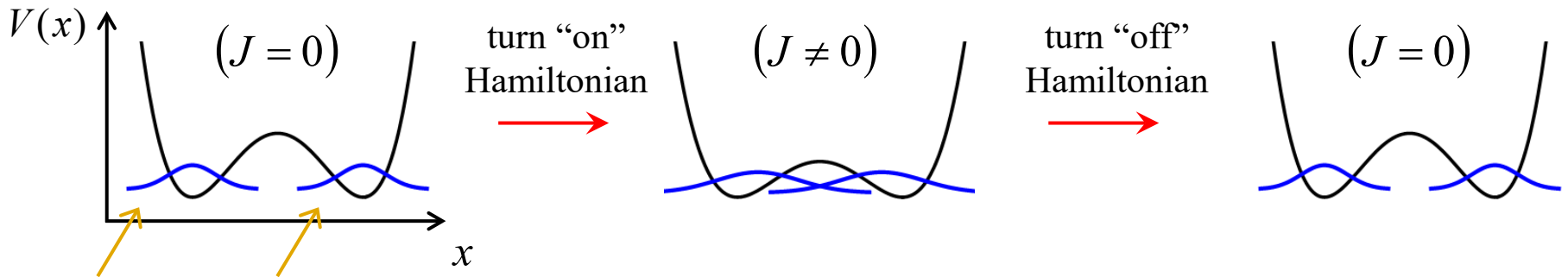
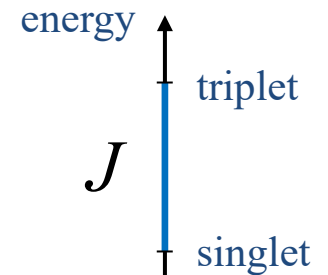
# Controlling Exchange

Petta *et al.*, *Science* (2005)



- Exchange Hamiltonian

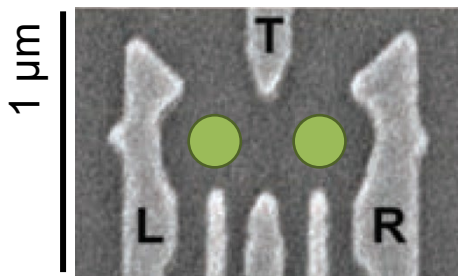
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Electron wave functions in quantum dot potential  $V(x)$

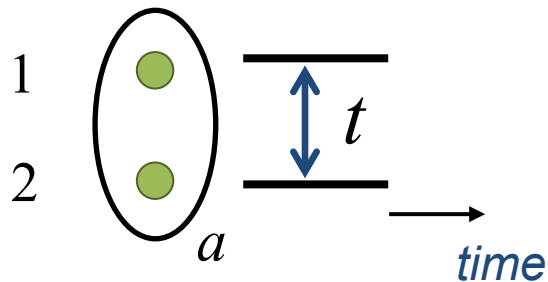
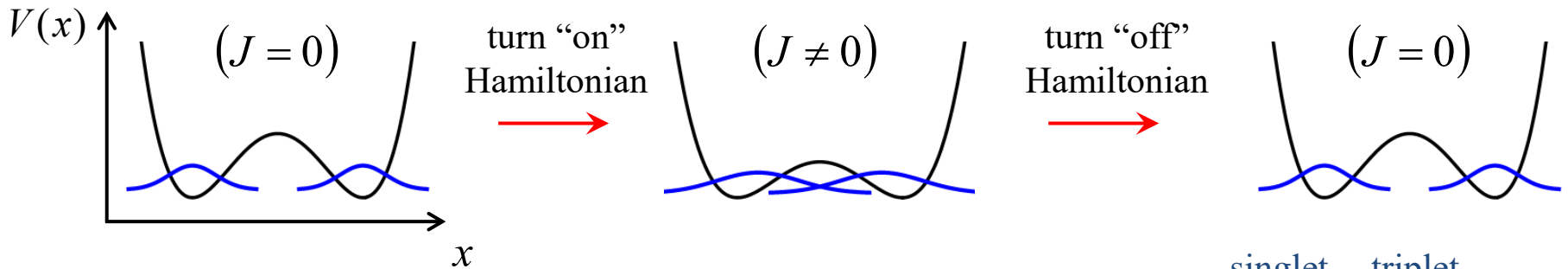
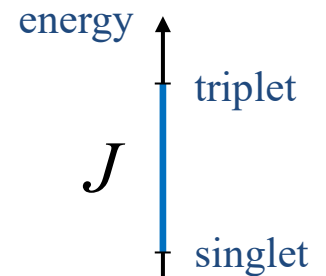
# Controlling Exchange

Petta *et al.*, *Science* (2005)



- Exchange Hamiltonian

$$H = J \mathbf{S}_1 \cdot \mathbf{S}_2$$

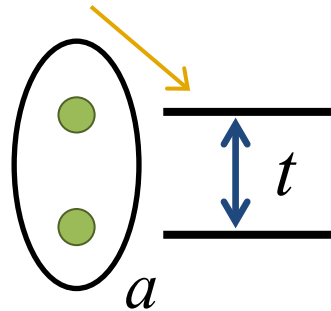


$a = 0$  ← singlet  
 $a = 1$  ← triplet

$$\exp(-iHt) = \begin{pmatrix} 1 & \\ & e^{-it} \end{pmatrix} \quad (J=1)$$

# Simple Exchange Pulses

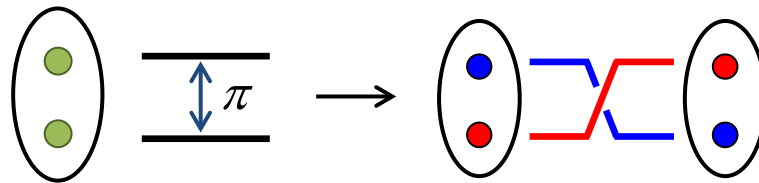
exchange pulse of duration  $t$



$$a = \begin{pmatrix} 0 & 1 \\ 1 & e^{-it} \end{pmatrix}$$

- SWAP pulse

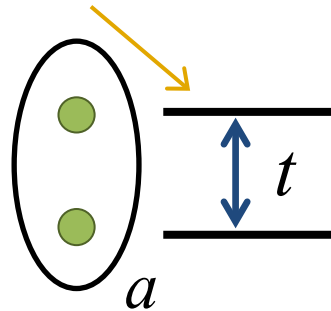
$$t = \pi$$



$$a = \begin{pmatrix} 0 & 1 \\ 1 & -1 \end{pmatrix}$$

# Simple Exchange Pulses

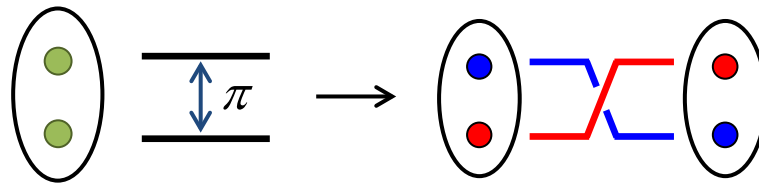
exchange pulse of duration  $t$



$$a = \begin{pmatrix} 0 & 1 \\ 1 & e^{-it} \end{pmatrix}$$

- SWAP pulse

$$t = \pi$$



$$a = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & -1 & -1 & 1 \end{pmatrix} = - \begin{pmatrix} -1 & 1 \\ 1 & 1 \end{pmatrix}$$

$$\text{singlet state} = |\uparrow\downarrow\rangle - |\downarrow\uparrow\rangle$$

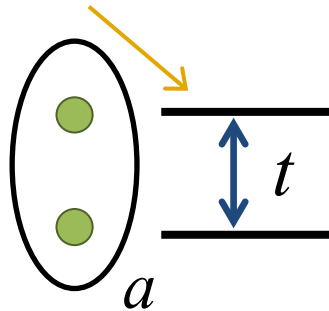
$(a = 0)$

$$\text{triplet states} = \begin{cases} |\uparrow\uparrow\rangle \\ |\uparrow\downarrow\rangle + |\downarrow\uparrow\rangle \\ |\downarrow\downarrow\rangle \end{cases}$$

$(a = 1)$

# Simple Exchange Pulses

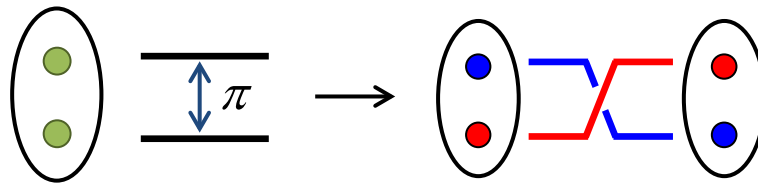
exchange pulse of duration  $t$



$$a = \begin{pmatrix} 0 & 1 \\ 1 & e^{-it} \end{pmatrix}$$

- SWAP pulse

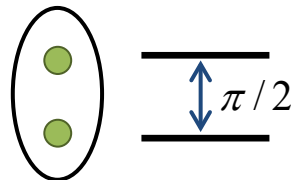
$$t = \pi$$



$$a = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & -1 & -1 & 1 \end{pmatrix} = - \begin{pmatrix} -1 & 1 \\ 1 & 1 \end{pmatrix}$$

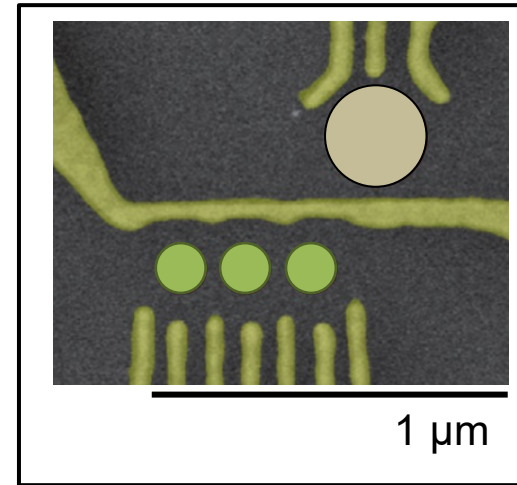
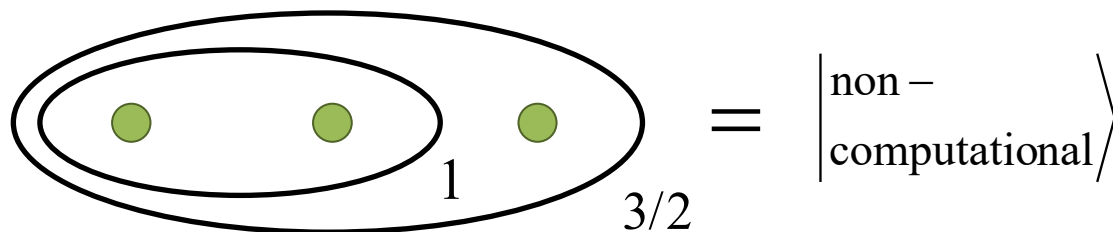
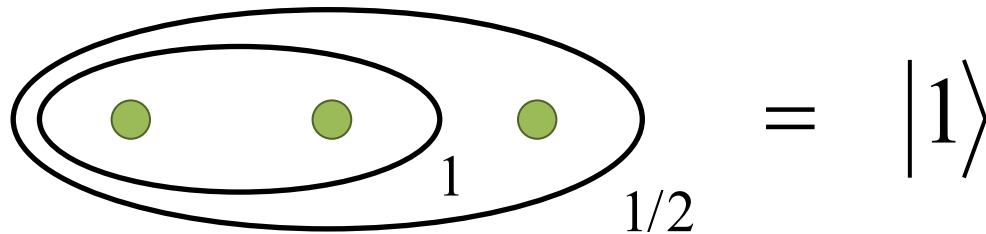
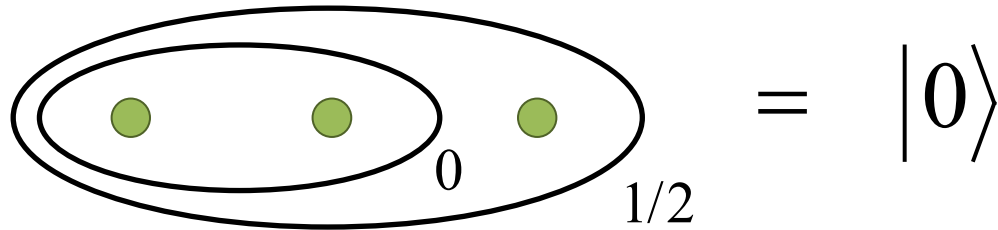
- SWAP<sup>1/2</sup> pulse

$$t = \pi/2$$

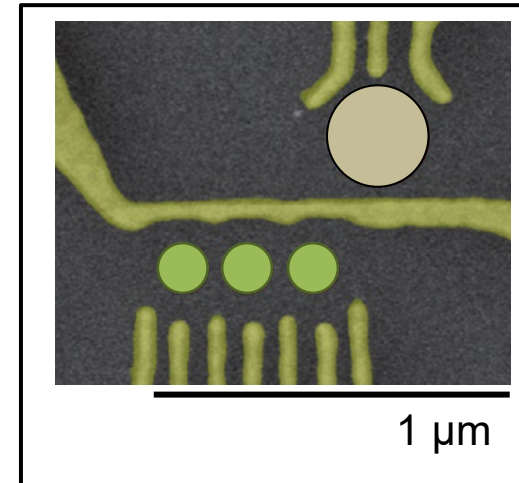
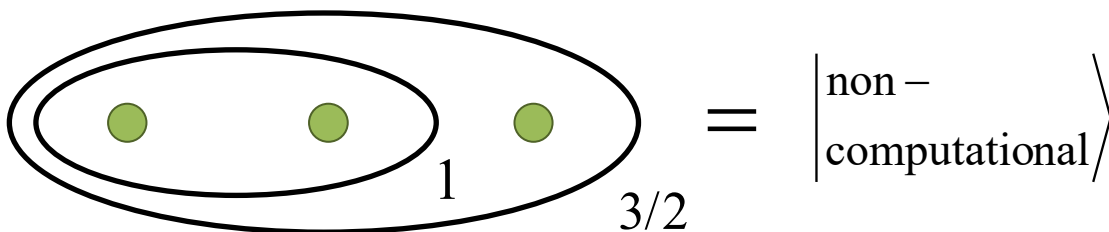
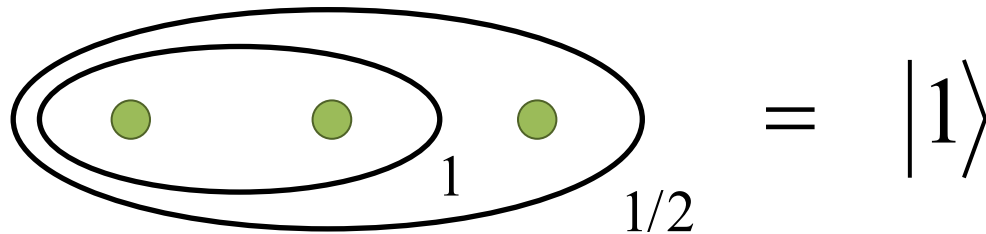
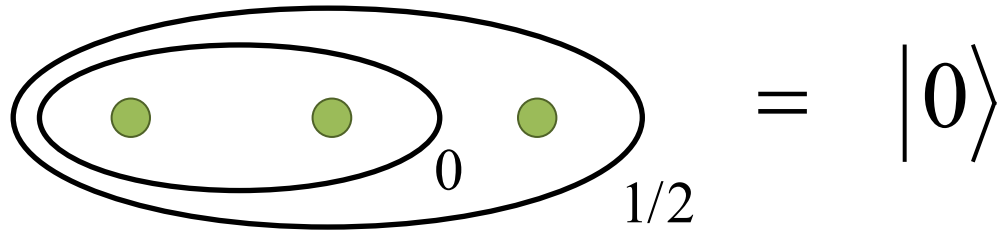




# Three-Spin Qubit Encoding



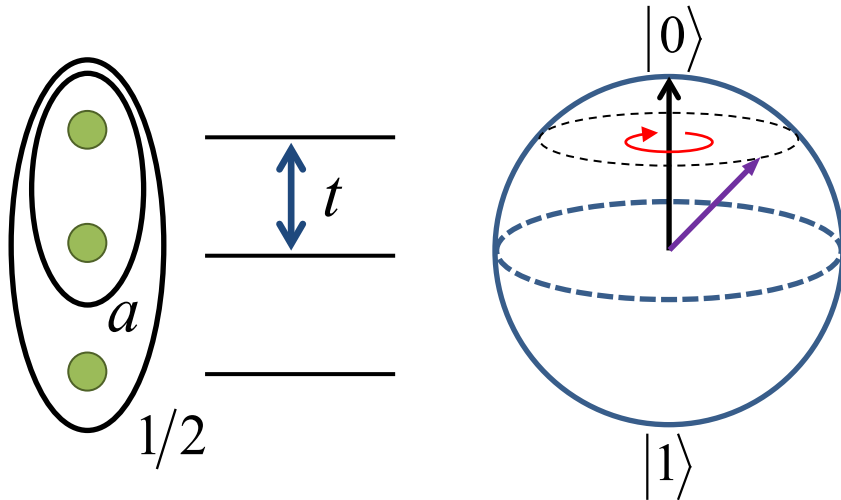
# Three-Spin Qubit Encoding



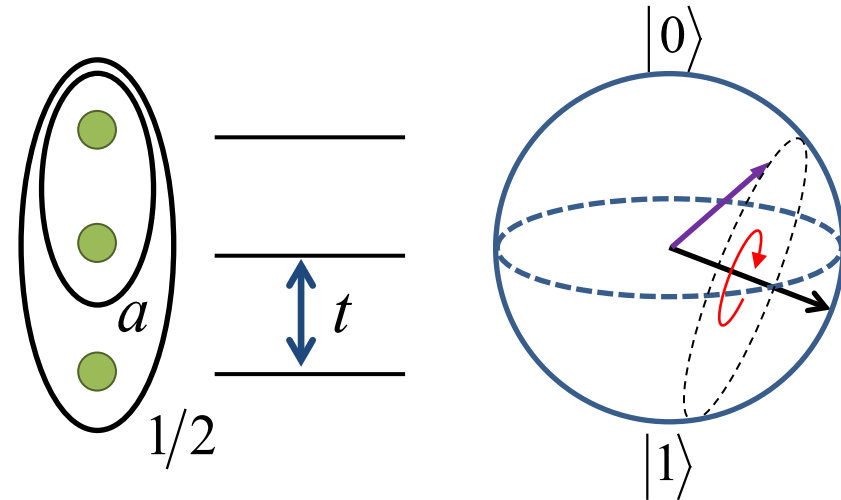
Transitions to this state are **leakage errors**.

# Single-Qubit Gates

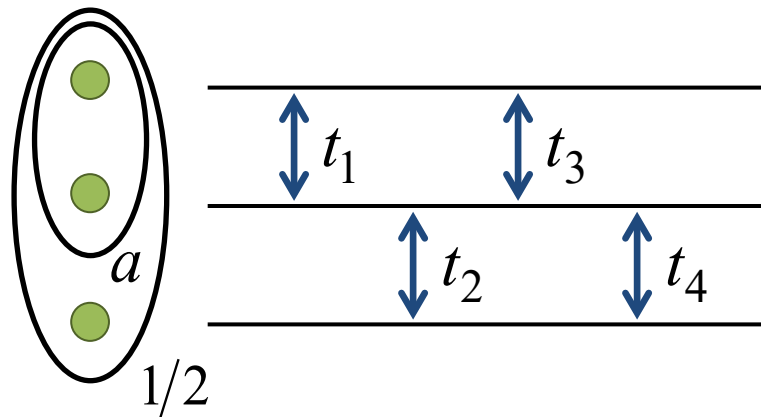
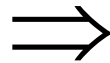
- Rotation about z-axis :



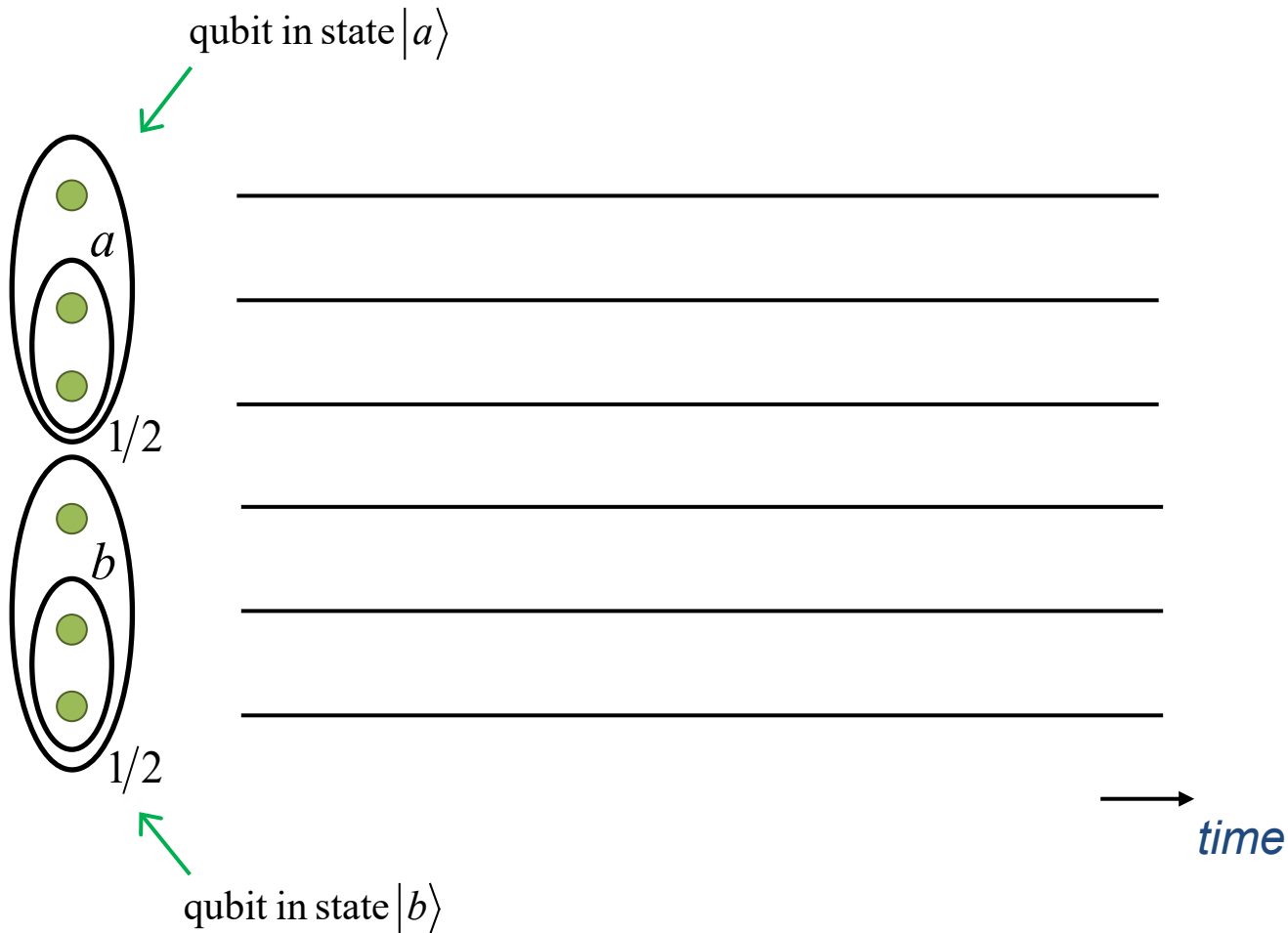
- Rotation about other axis:



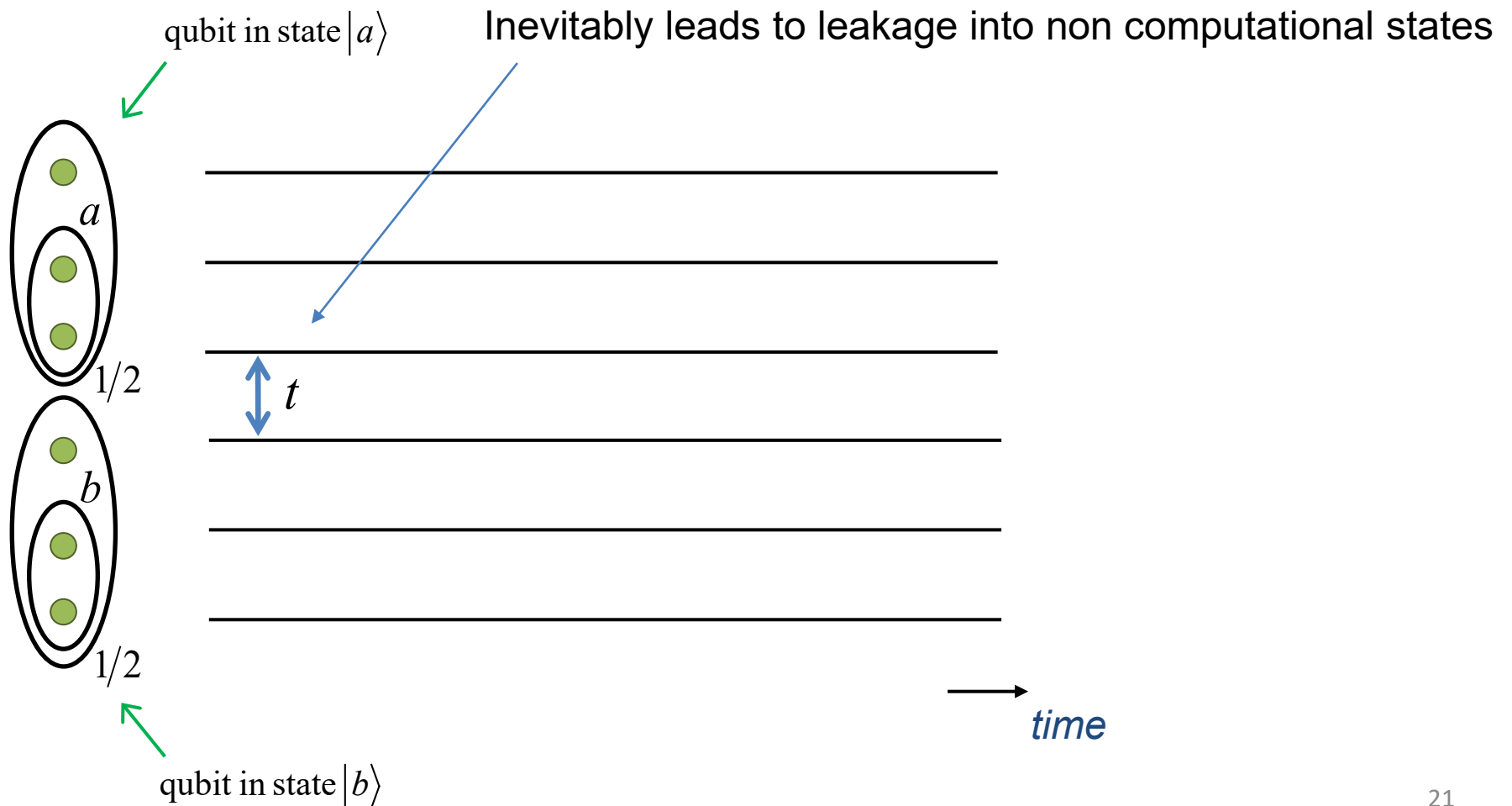
arbitrary  
rotations



# Two-Qubit Gates



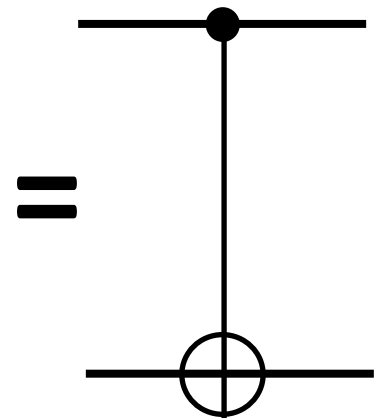
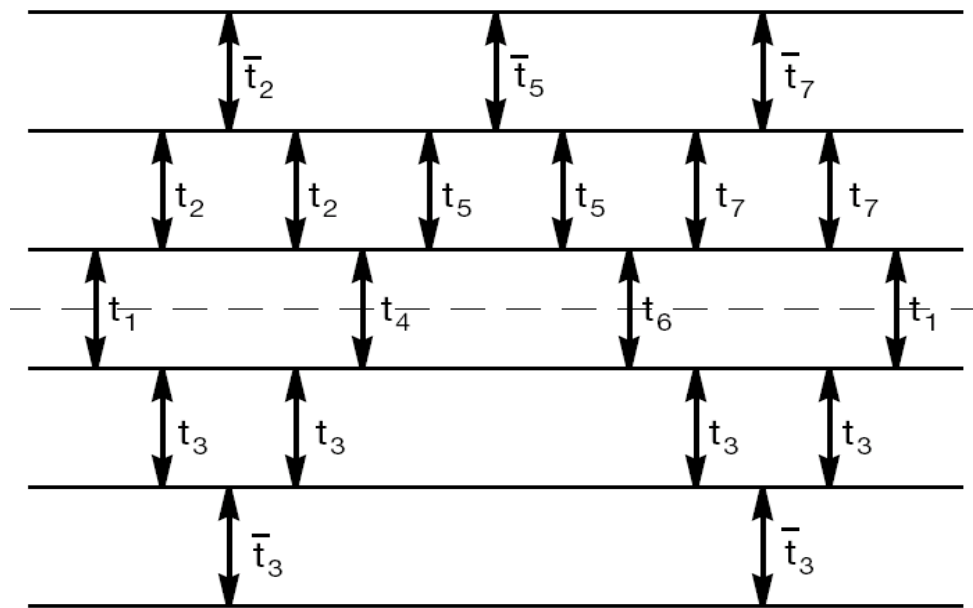
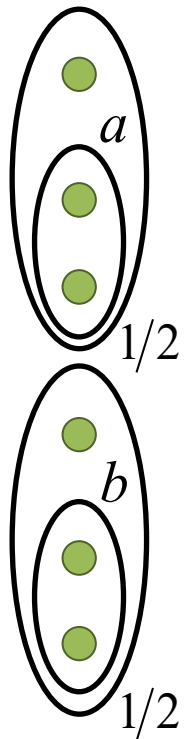
# Two-Qubit Gates



# Two-Qubit Gates

DiVincenzo, Bacon, Kempe, Burkard & Whaley, Nature (2000)

19 pulse sequence found numerically

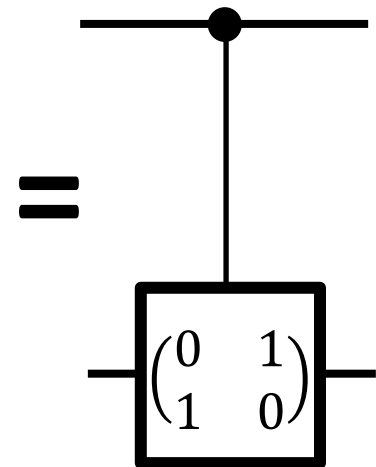
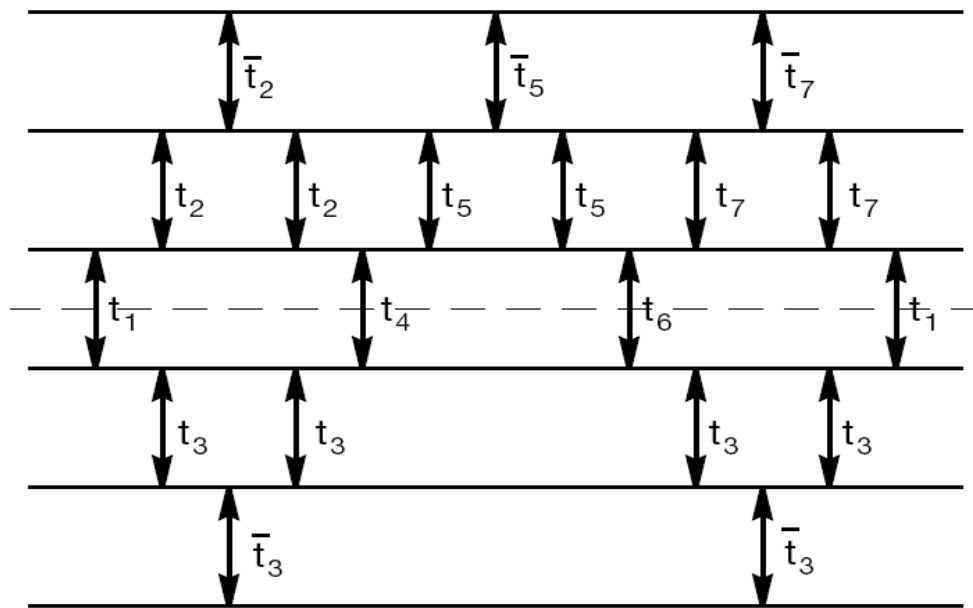
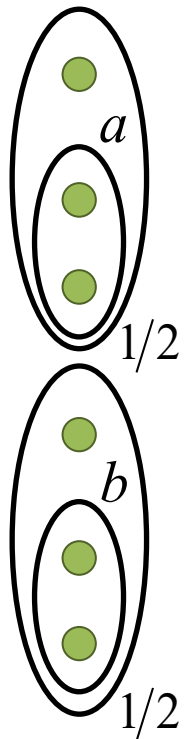


$$t_1 = 2.581\dots, t_2 = 1.303\dots, t_3 = 1.753\dots, \dots$$

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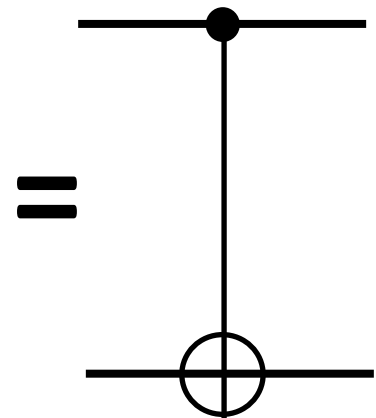
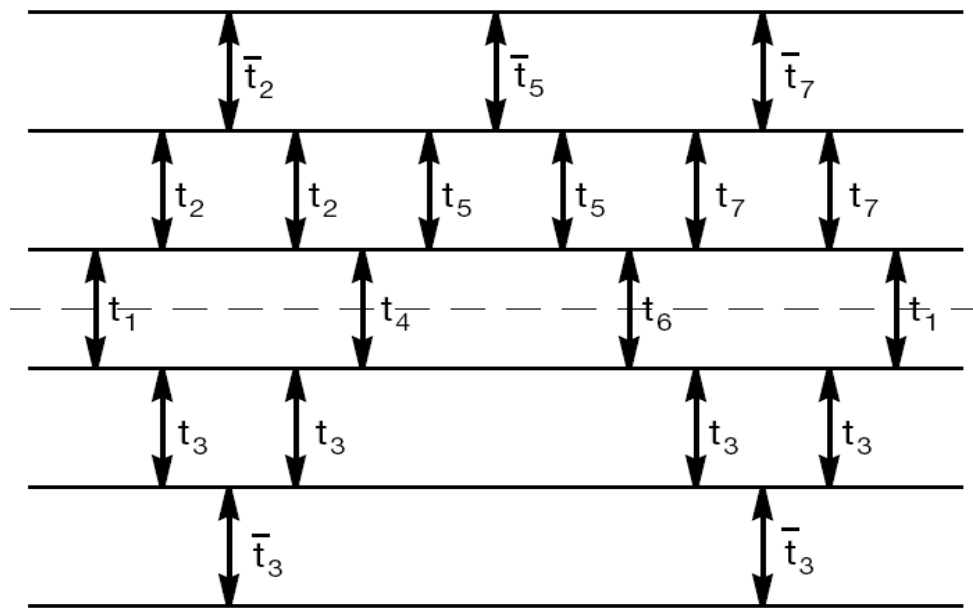
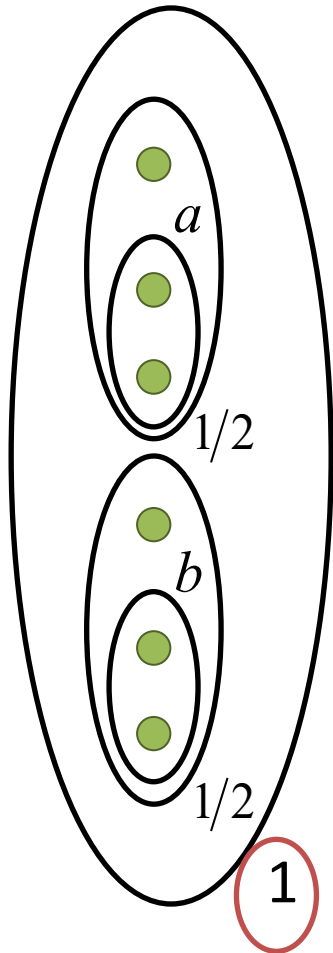


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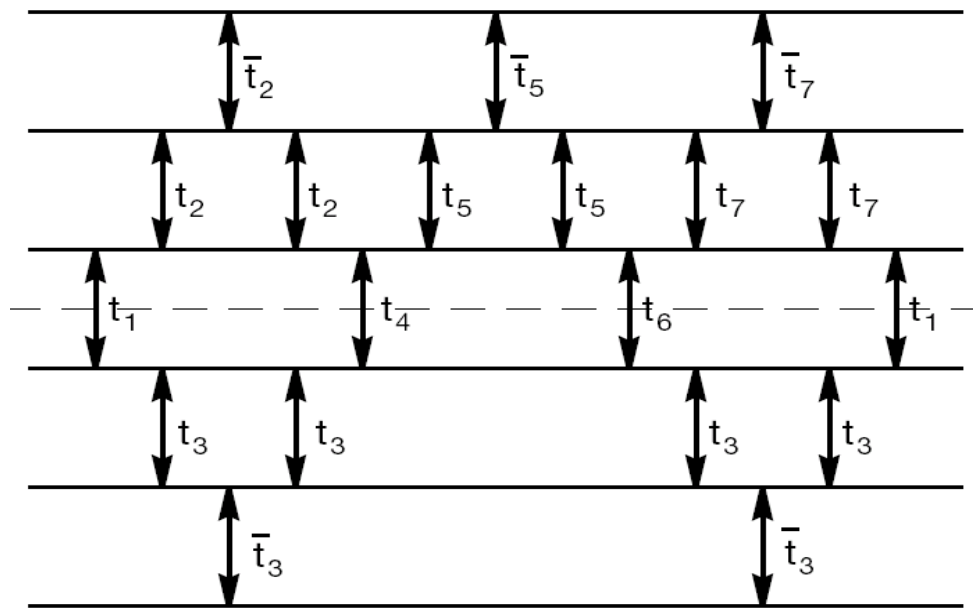
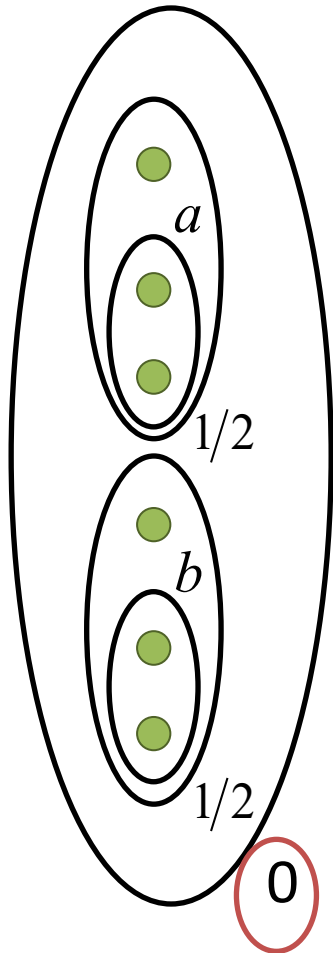
← Gives a CNOT (up to single qubit operations) if the total spin is 1



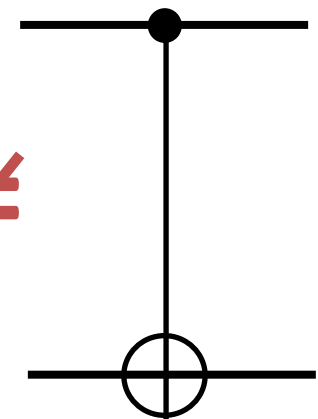
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DiVincenzo, Bacon, Kempe, Burkard & Whaley, Nature (2000)

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$\neq$

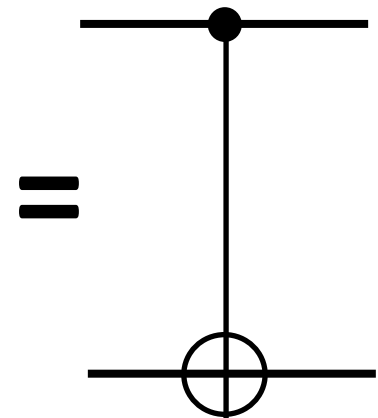
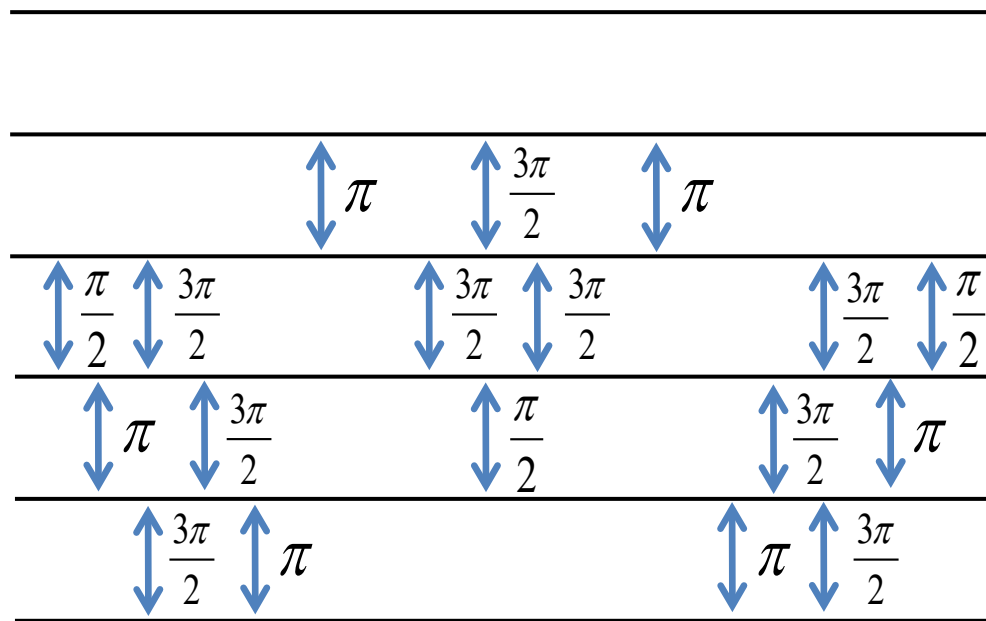
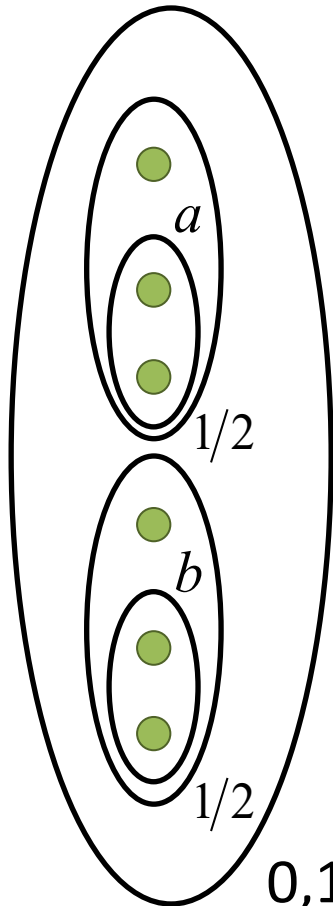


← Does **not** give a CNOT if the total spin is  $0$

# Fong-Wandzura Sequence

Fong & Wandzura, *Quantum Information and Computation* (2011)

18 pulse sequence found numerically

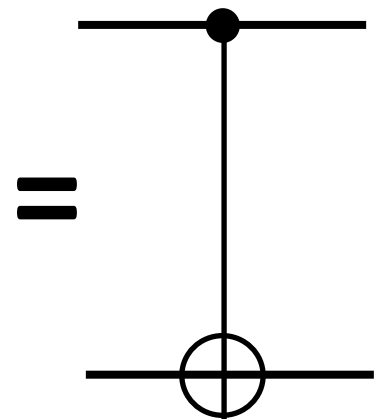
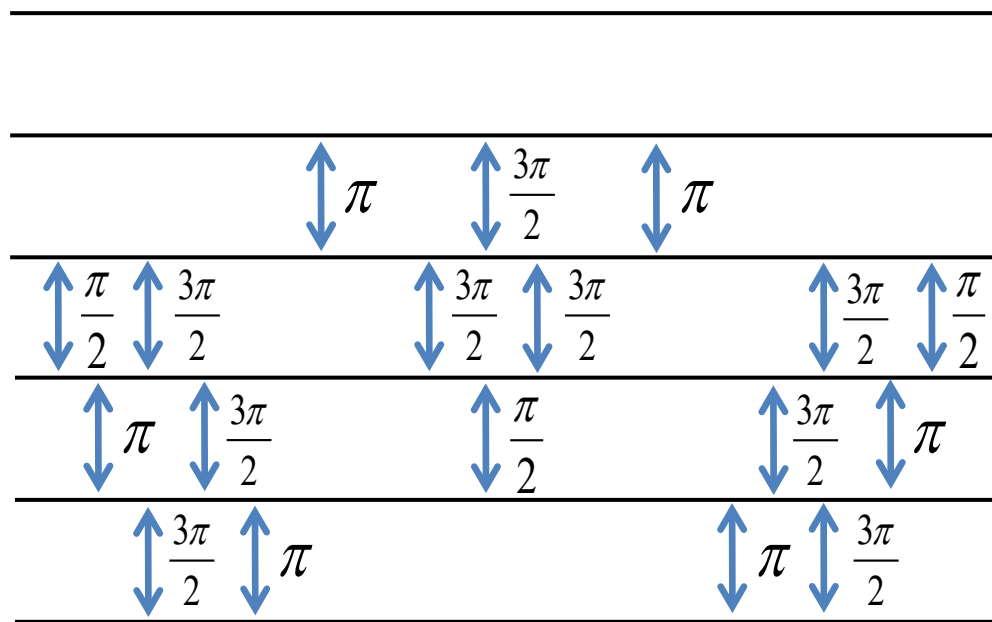
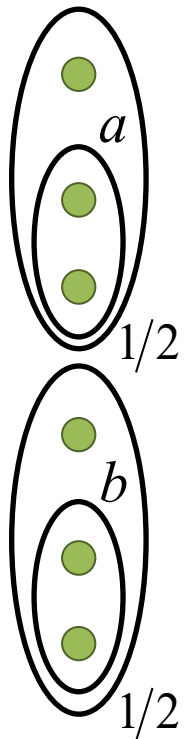


0,1 ← Gives CNOT for *both* total spin 0 and 1

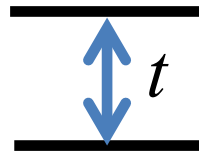
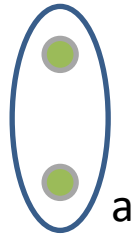
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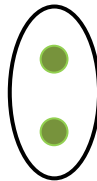
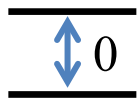
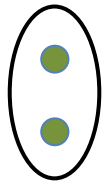


# Two Simple Pulses



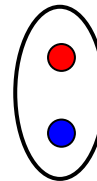
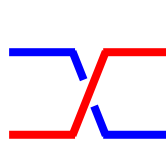
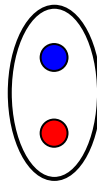
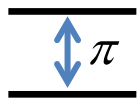
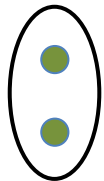
$$a = \begin{pmatrix} 0 & 1 \\ 1 & e^{-it} \end{pmatrix}$$

$t = 0$



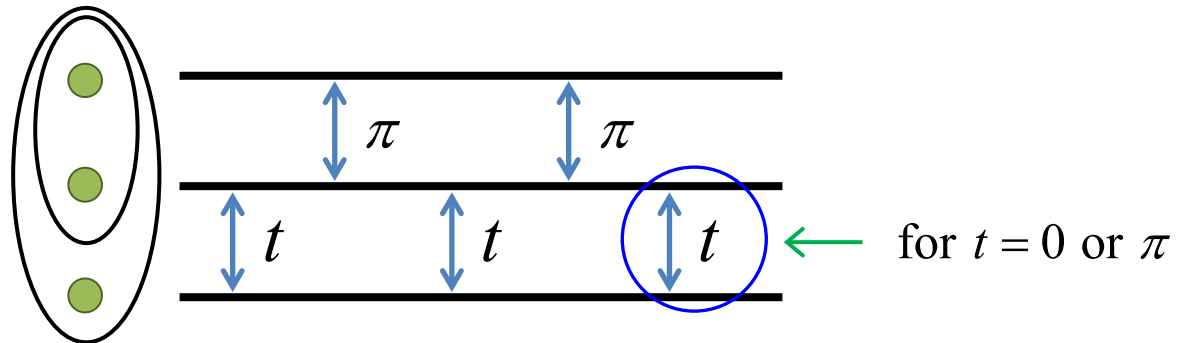
$$\begin{pmatrix} 1 & \\ & 1 \end{pmatrix}$$

$t = \pi$

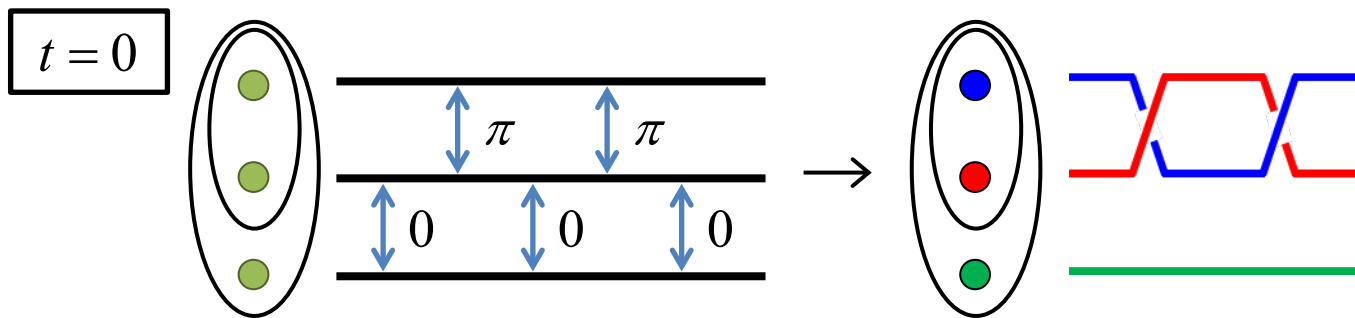
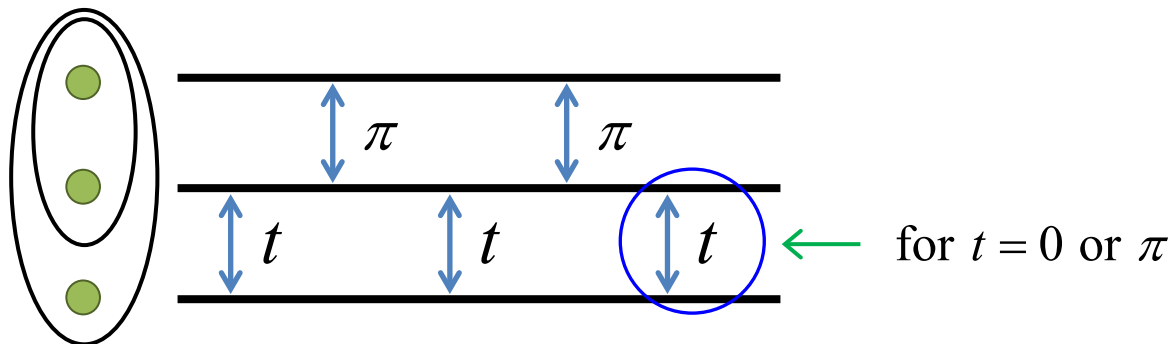


$$\begin{pmatrix} 1 & \\ & -1 \end{pmatrix}$$

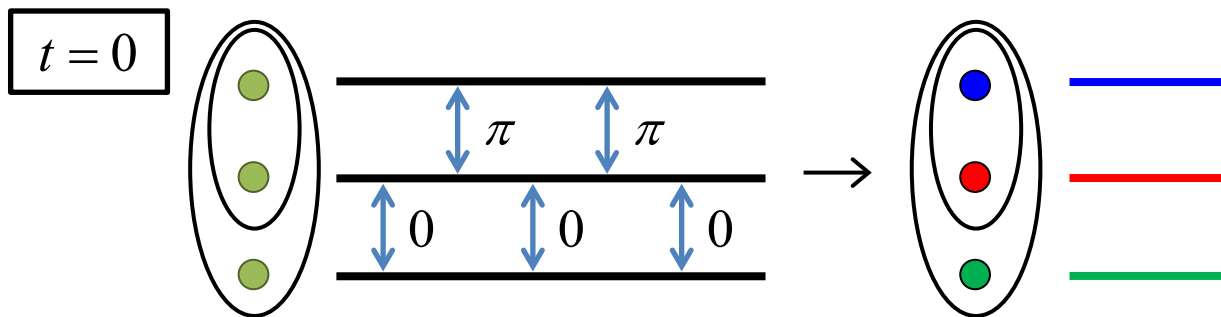
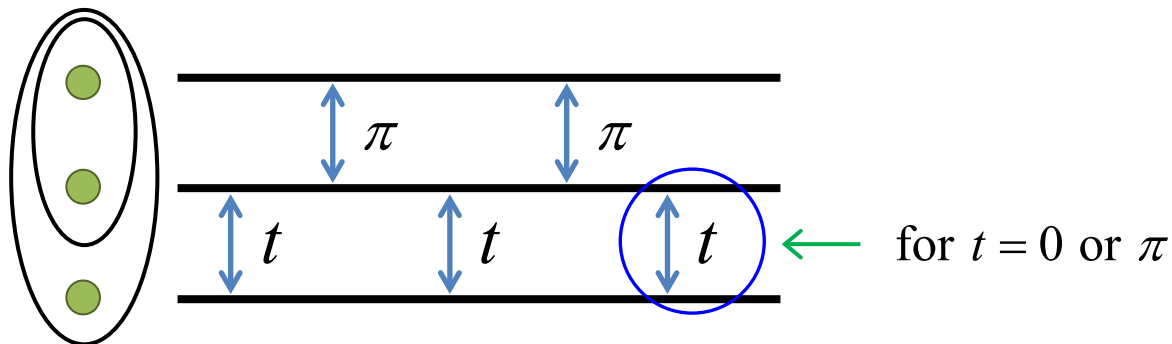
# Two Simple Sequences



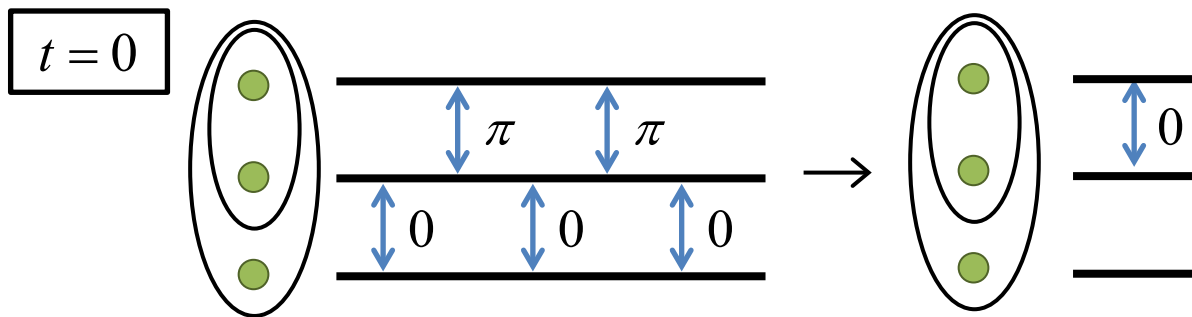
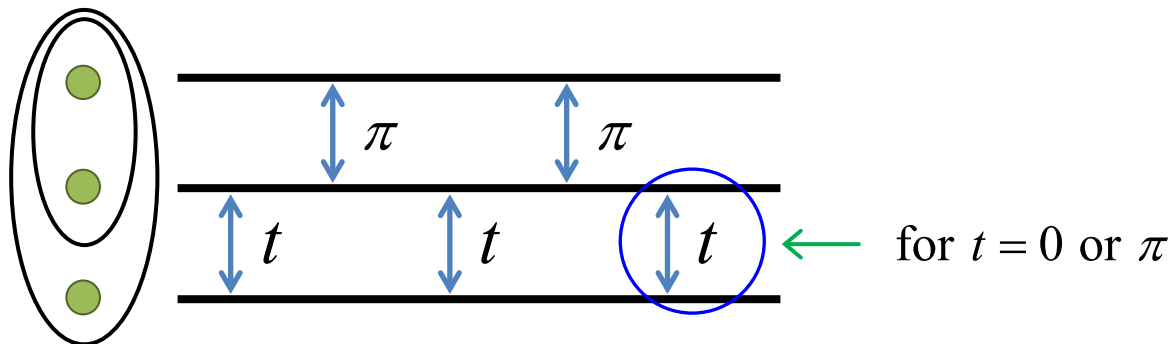
# Two Simple Sequences



# Two Simple Sequences

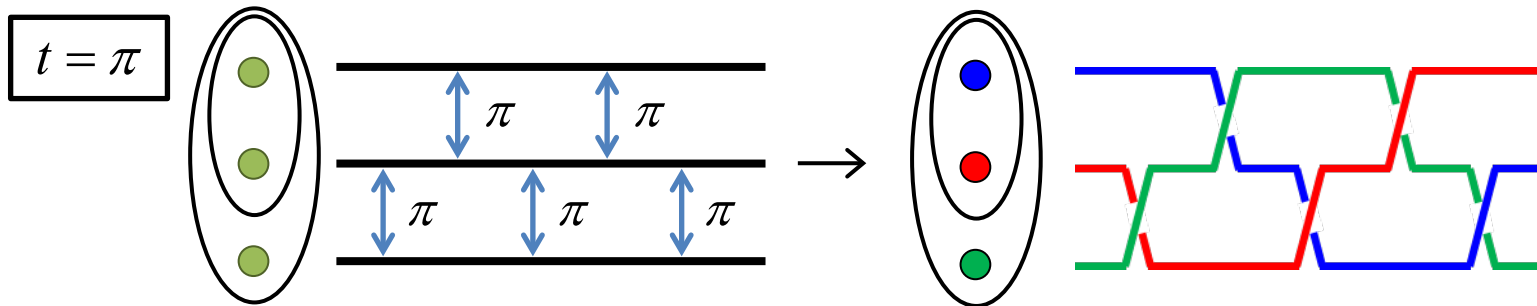
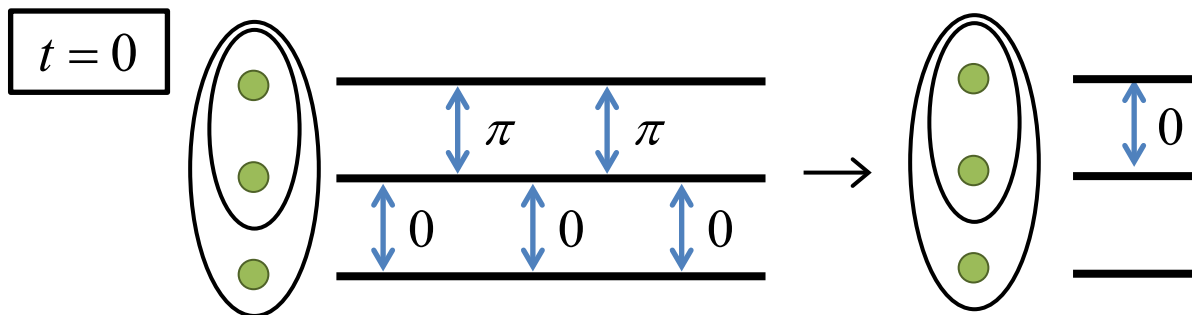
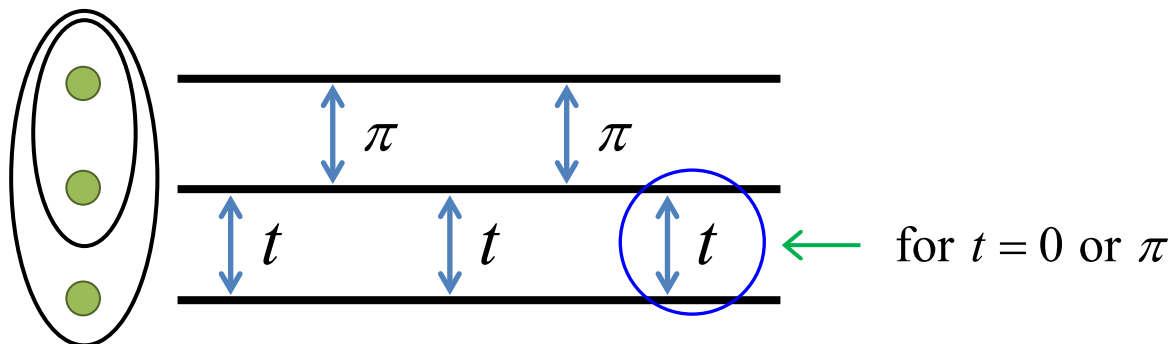


# Two Simple Sequences

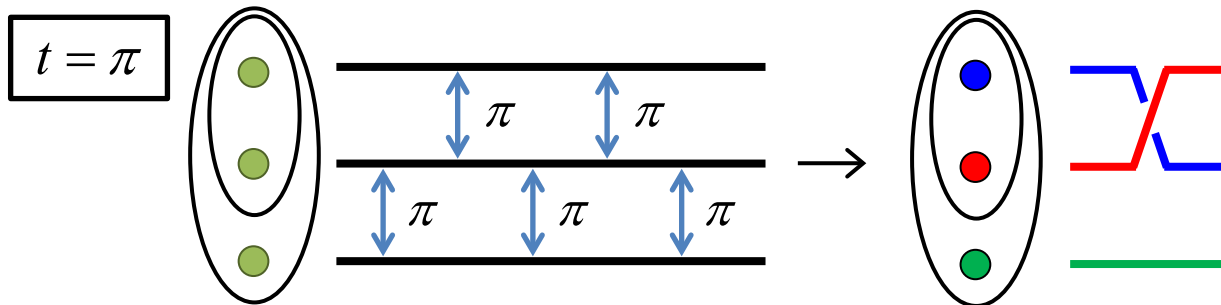
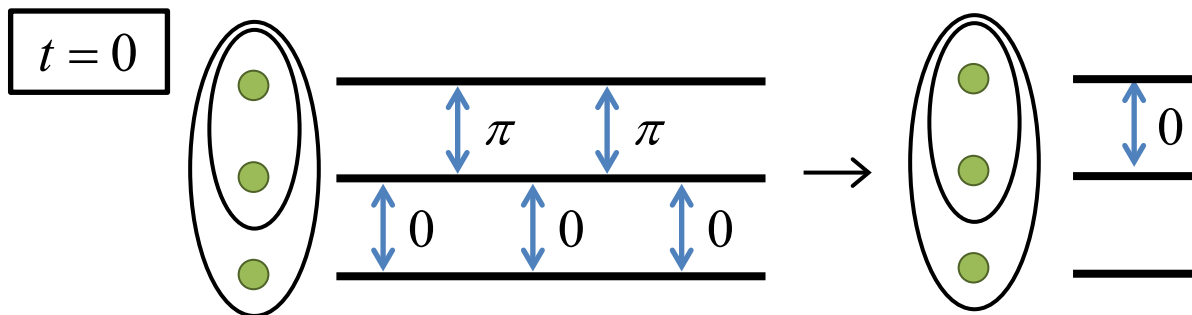
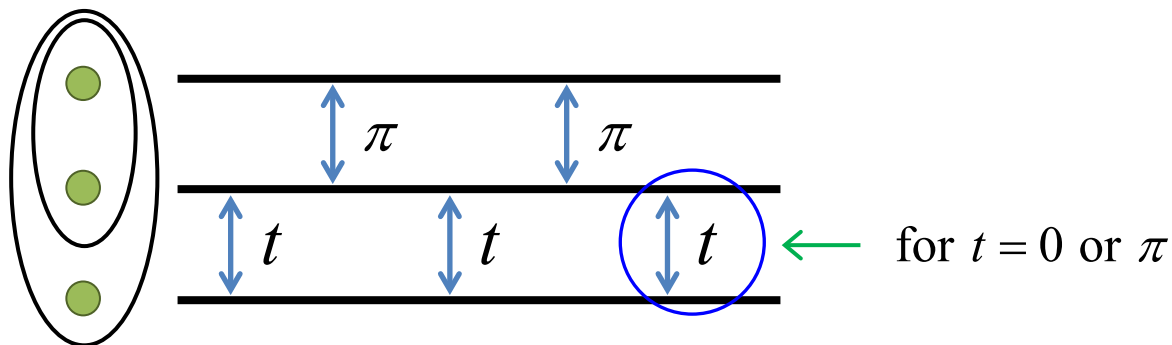




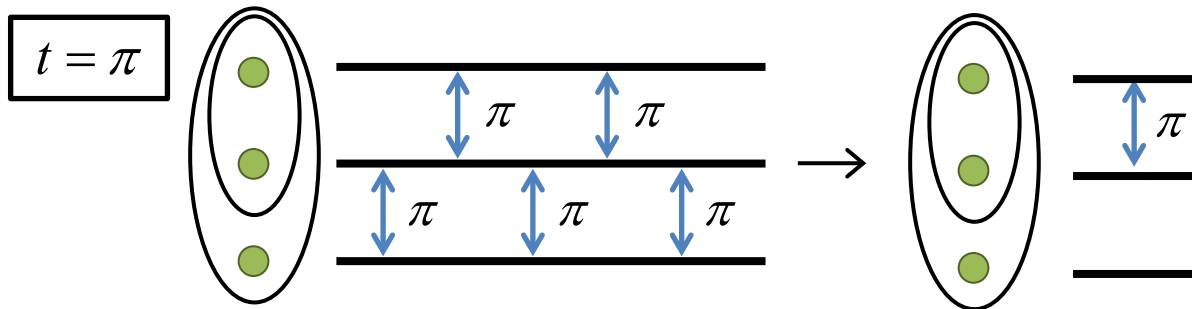
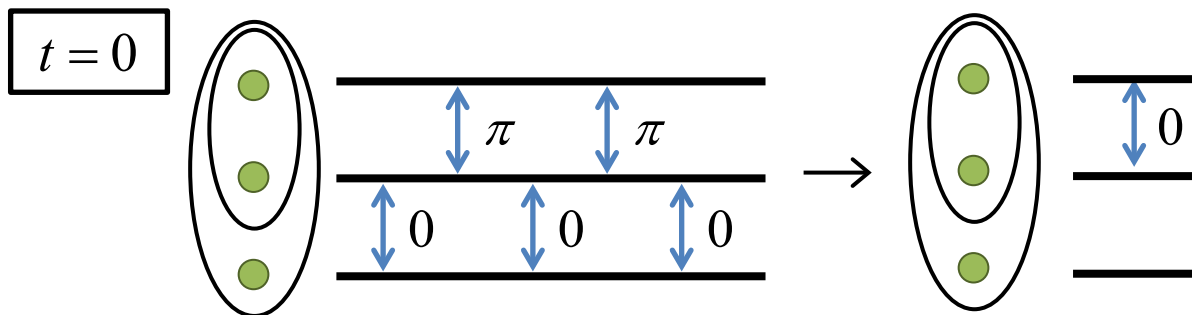
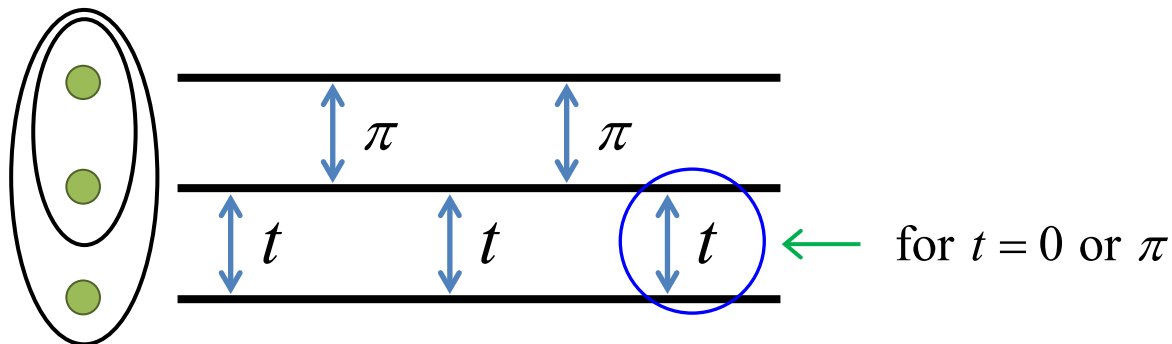
# Two Simple Sequences



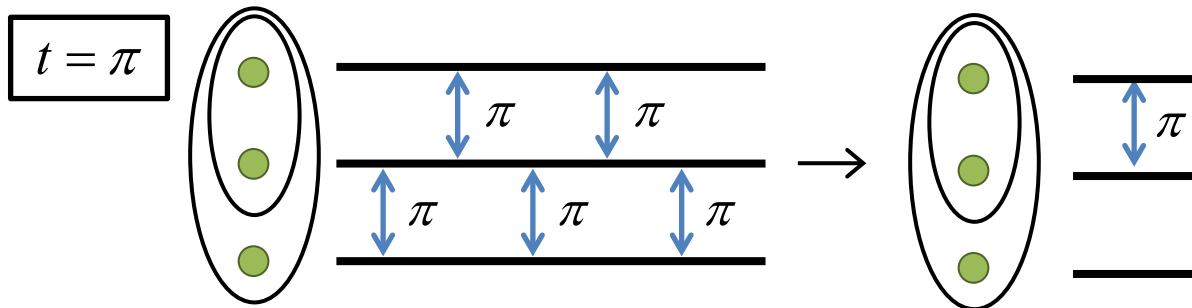
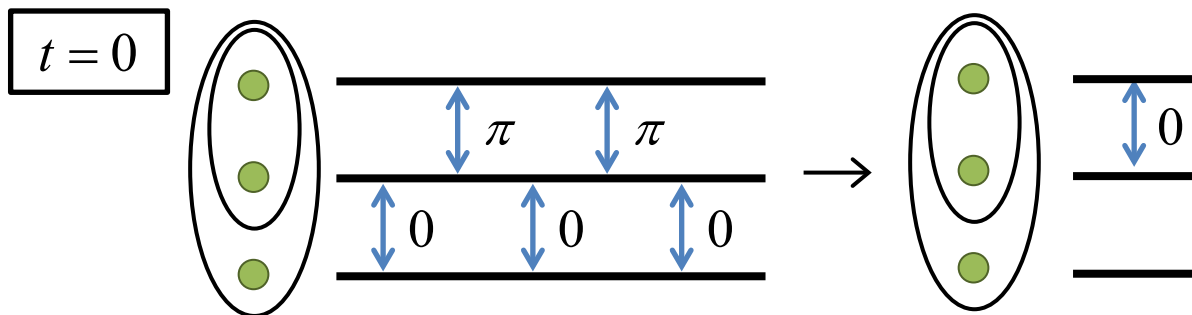
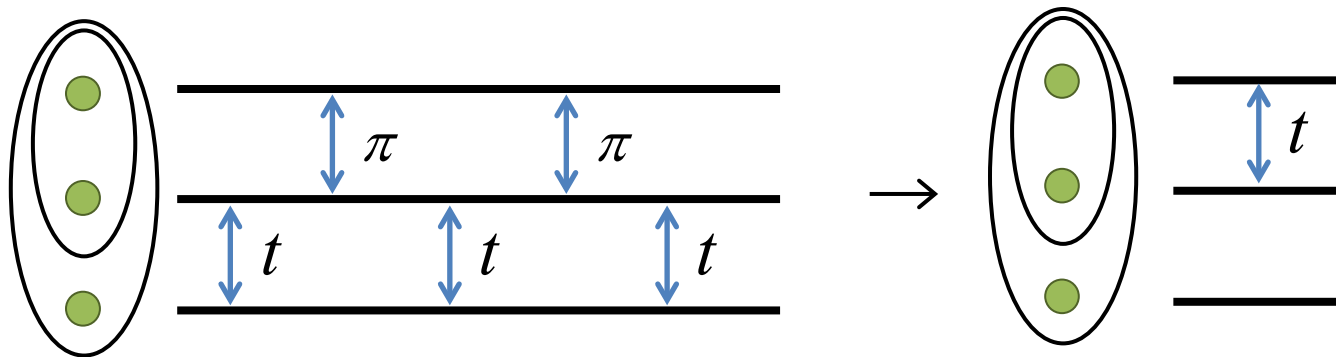
# Two Simple Sequences



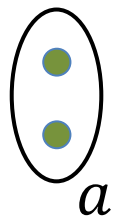
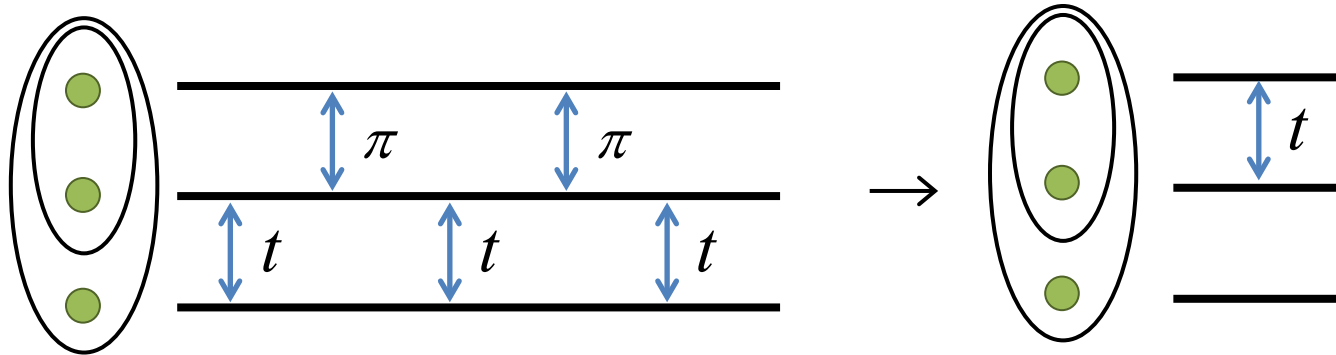
# Two Simple Sequences



# Two Simple Sequences



# Sequence Identity



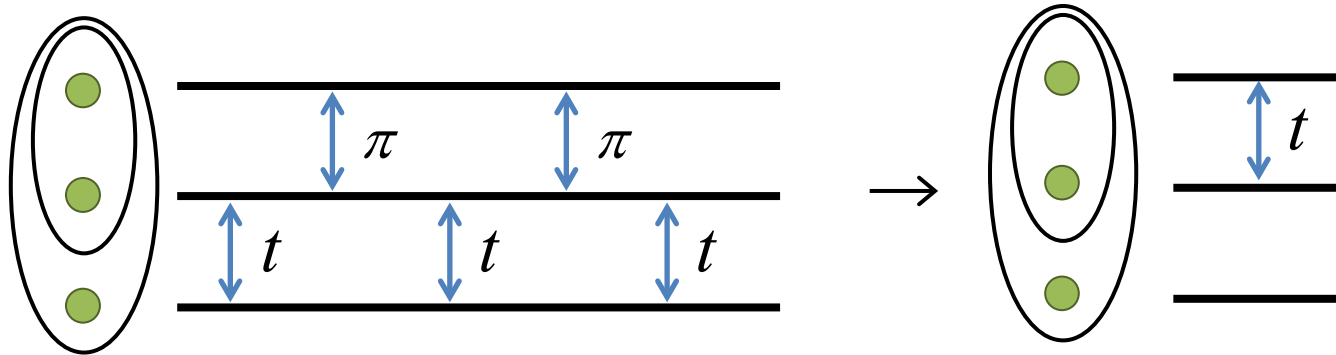
$$a = \begin{pmatrix} 0 & 1 \\ 1 & m \end{pmatrix}$$

Identity holds for

$$\begin{aligned} t = 0 &\rightarrow m = +1 \\ t = \pi &\rightarrow m = -1 \end{aligned}$$

$$m^2 = 1$$

# Sequence Identity



$$\begin{array}{c}
 a = \begin{matrix} 0 & 1 \\ \left( \begin{matrix} 1 & \\ & m \end{matrix} \right) \end{matrix} \\
 \begin{array}{c} \text{Oval} \\ \text{with 2 dots} \\ \text{distance } t \end{array}
 \end{array}$$

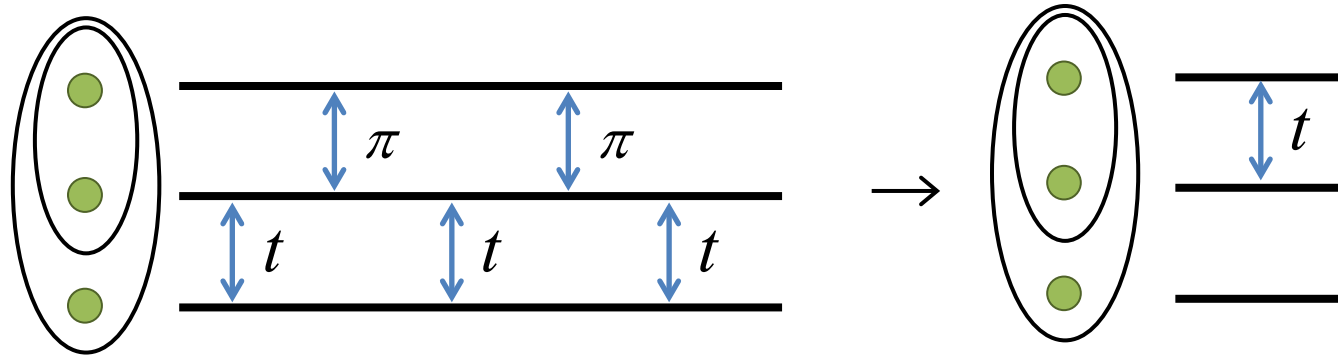
Identity holds for

$$\begin{array}{l}
 t = 0 \rightarrow m = +1 \\
 t = \pi \rightarrow m = -1
 \end{array}$$

$$\boxed{m^2 = 1}$$

Can show: Any pulse with  $m^2=1$  will satisfy this identity

# Sequence Identity



$$\begin{array}{c}
 a = \begin{matrix} 0 & 1 \\ \begin{pmatrix} 1 & \\ & m \end{pmatrix} \end{matrix} \\
 \begin{array}{c} \text{Oval with 2 dots} \\ \text{---} \\ \updownarrow t \\ \text{---} \end{array} \\
 a
 \end{array}$$

Identity holds for

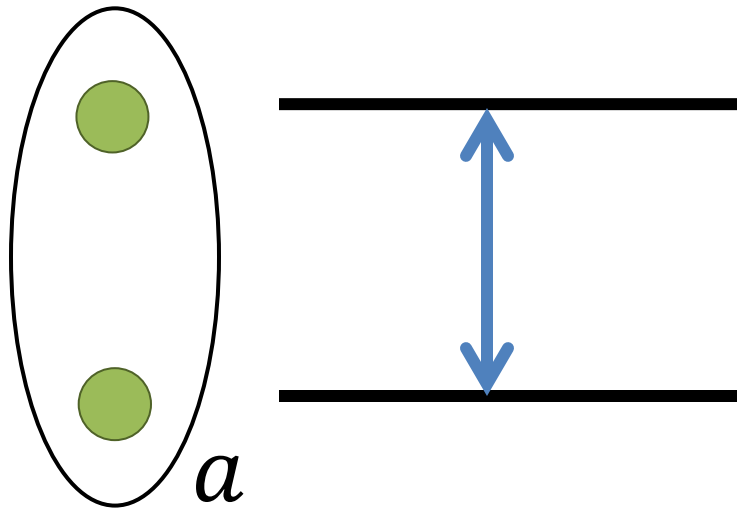
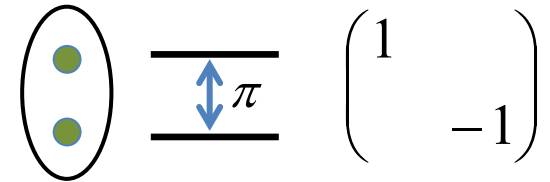
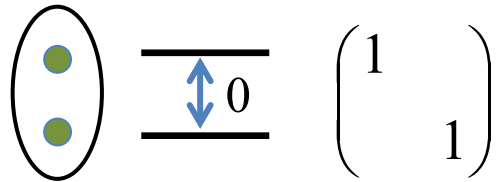
$$\begin{array}{l}
 t = 0 \rightarrow m = +1 \\
 t = \pi \rightarrow m = -1
 \end{array}$$

$$m^2 = 1$$

Can show: Any pulse with  $m^2=1$  will satisfy this identity

Does  $m$  have to be a number?  
How about a matrix?

# $m^2 = 1$ Pulses

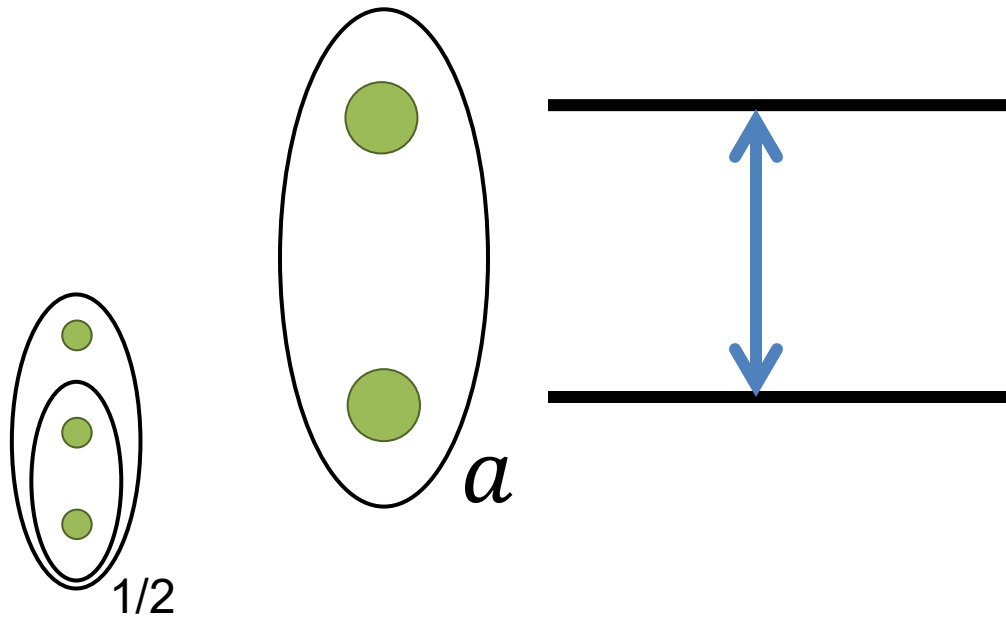
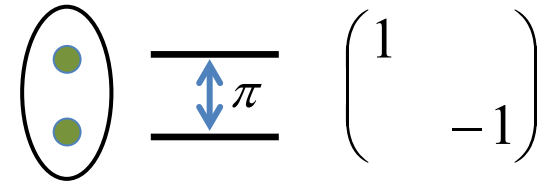
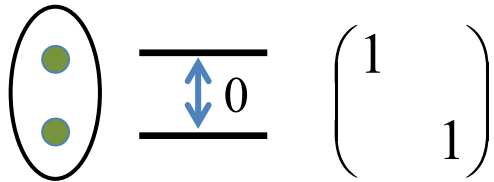


$$a = \begin{pmatrix} 0 & 1 \\ 1 & m \end{pmatrix}$$

$$m^2 = 1 \longrightarrow m = +1, -1$$



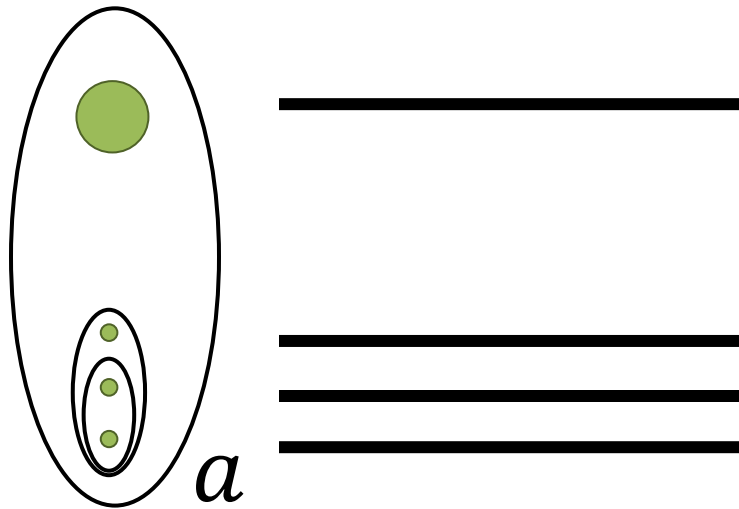
# $m^2 = 1$ Pulses



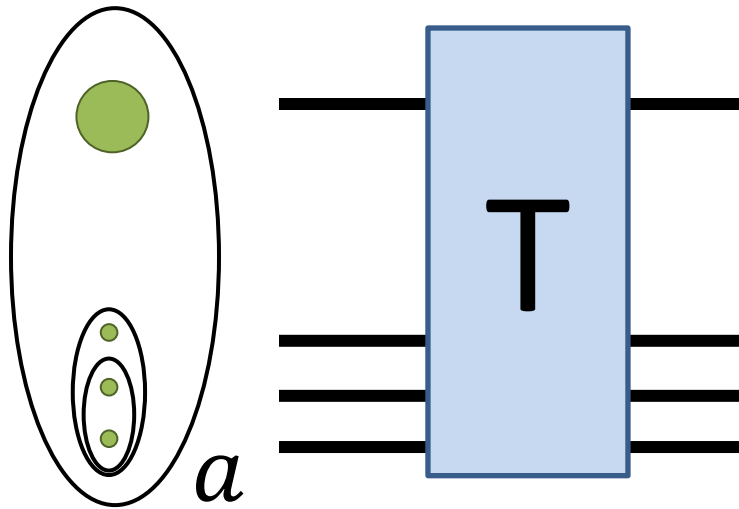
$$a = \begin{pmatrix} 0 & 1 \\ 1 & m \end{pmatrix}$$

$$m^2 = 1 \longrightarrow m = +1, -1$$

# “Elevating” $m^2 = 1$ Pulses

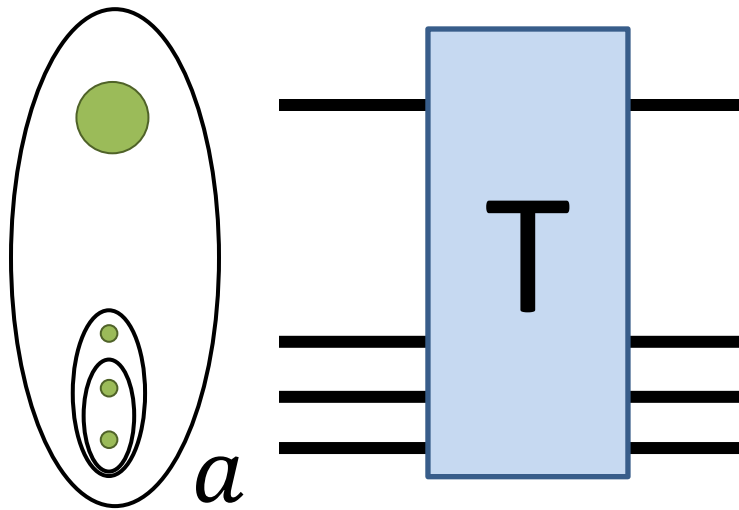


# “Elevating” $m^2 = 1$ Pulses



$$a = \begin{matrix} 0 & 1 \\ \left( \begin{matrix} I & \\ & M \end{matrix} \right) \end{matrix}$$

# “Elevating” $m^2 = 1$ Pulses

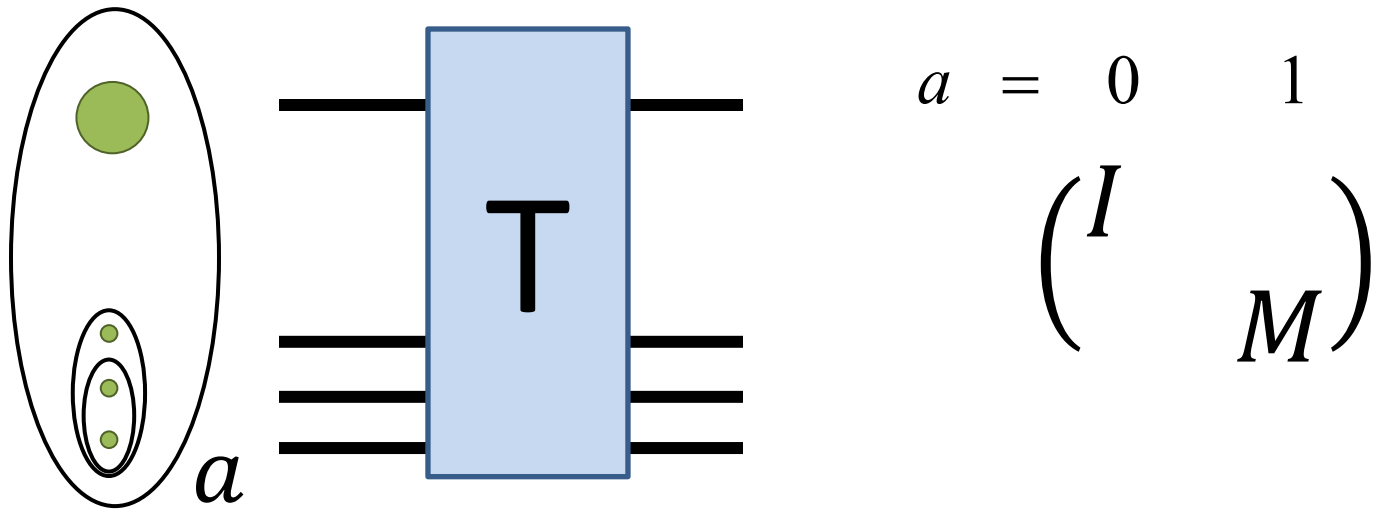


$$a = \begin{matrix} 0 & 1 \\ \left( \begin{matrix} I & \\ & M \end{matrix} \right) \end{matrix}$$

The diagram shows a 2x2 matrix structure. The top row is labeled 'a = 0' and '1'. The matrix elements are 'I' and 'M', each enclosed in a dashed circle. Blue arrows point from the text below to the 'I' and 'M' elements.

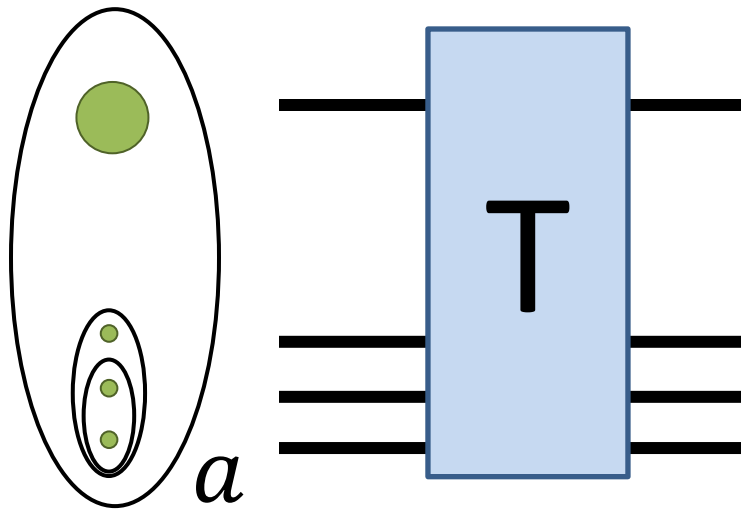
$2 \times 2$  matrices, acting on “swapped in” qubit

$$M^2 = I \text{ "Pulse"}$$



$$M^2 = I \rightarrow M = +I, -I$$

$$M^2 = I \text{ "Pulse"}$$



$$a = \begin{pmatrix} 0 & 1 \\ I & M \end{pmatrix}$$

$$M^2 = I \rightarrow M = +I, -I$$

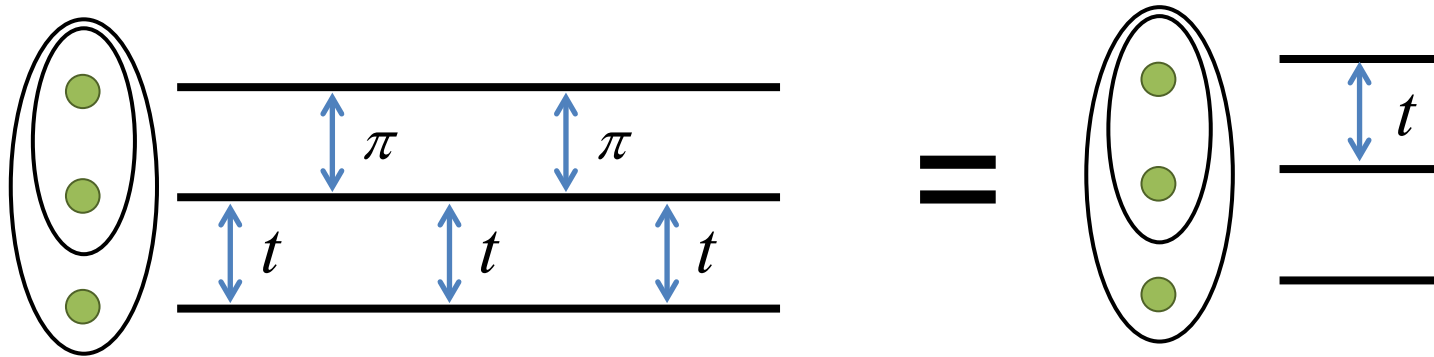
$$M = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

More solutions!

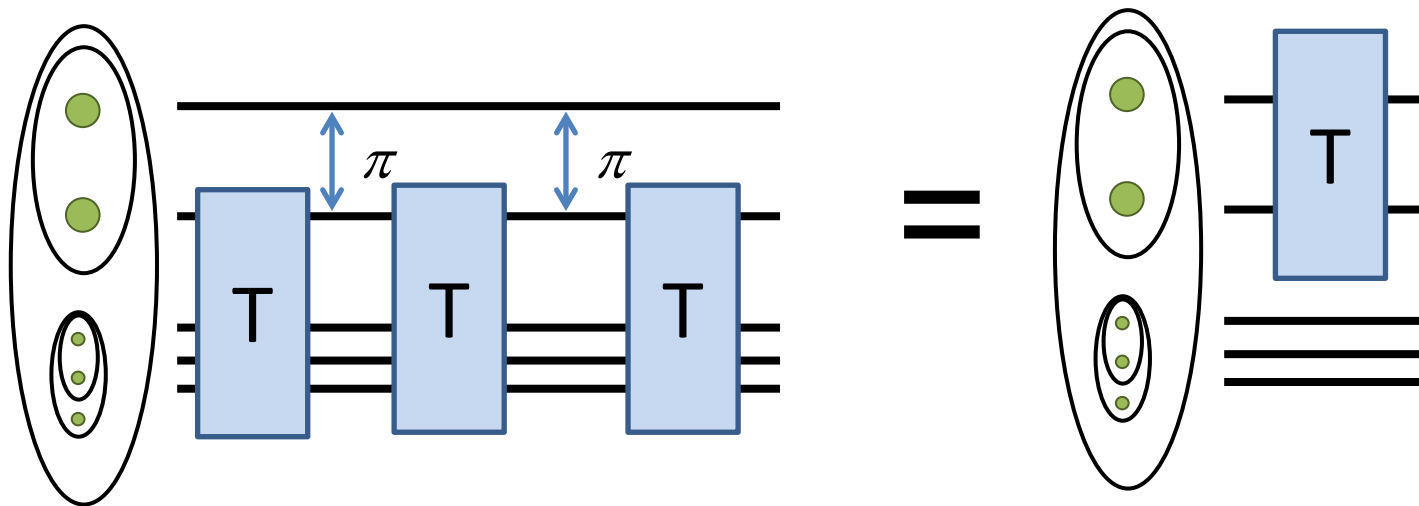
$$\rightarrow M = \hat{n} \cdot \sigma$$

$$M = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \dots$$

# Sequence Elevation

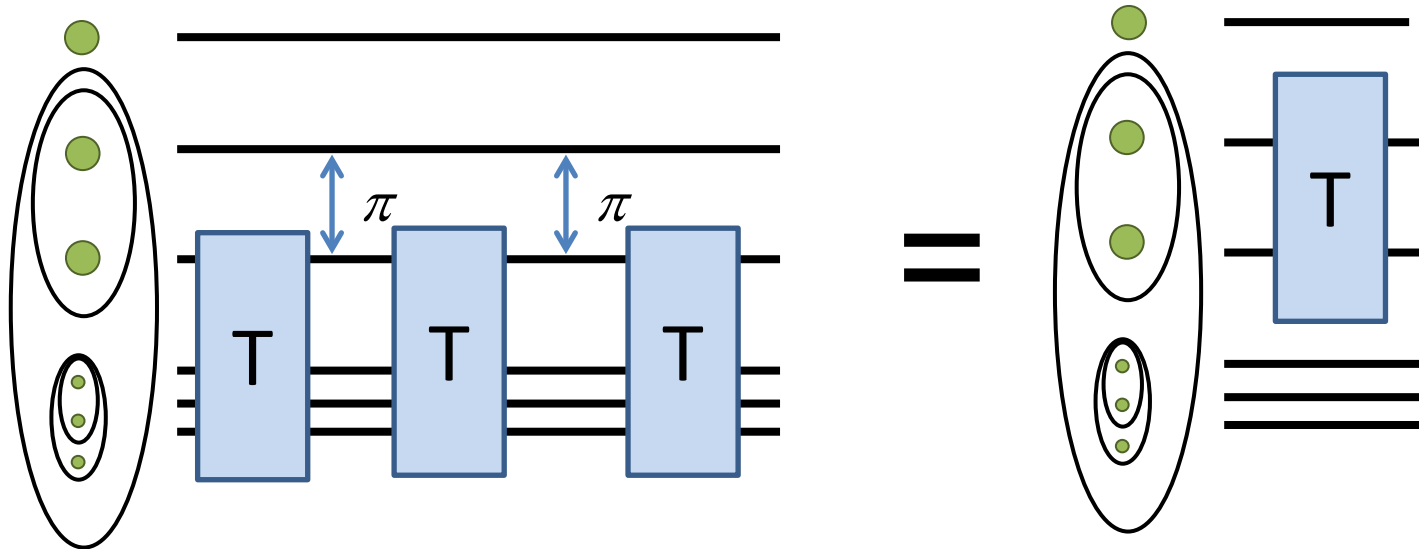


# Sequence Elevation

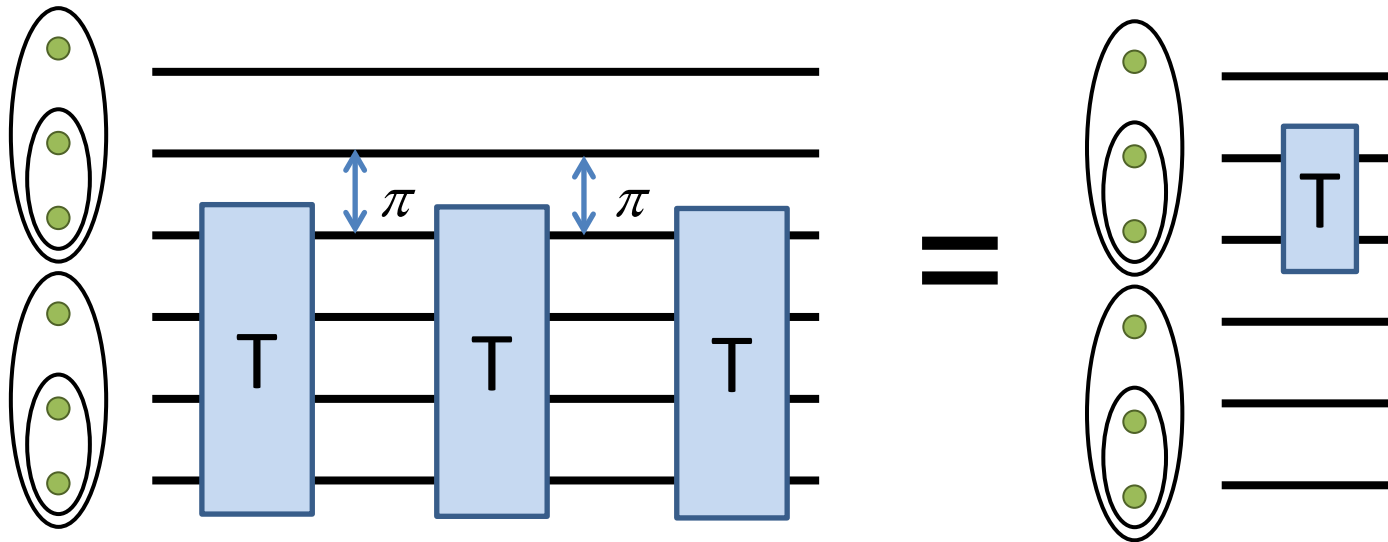




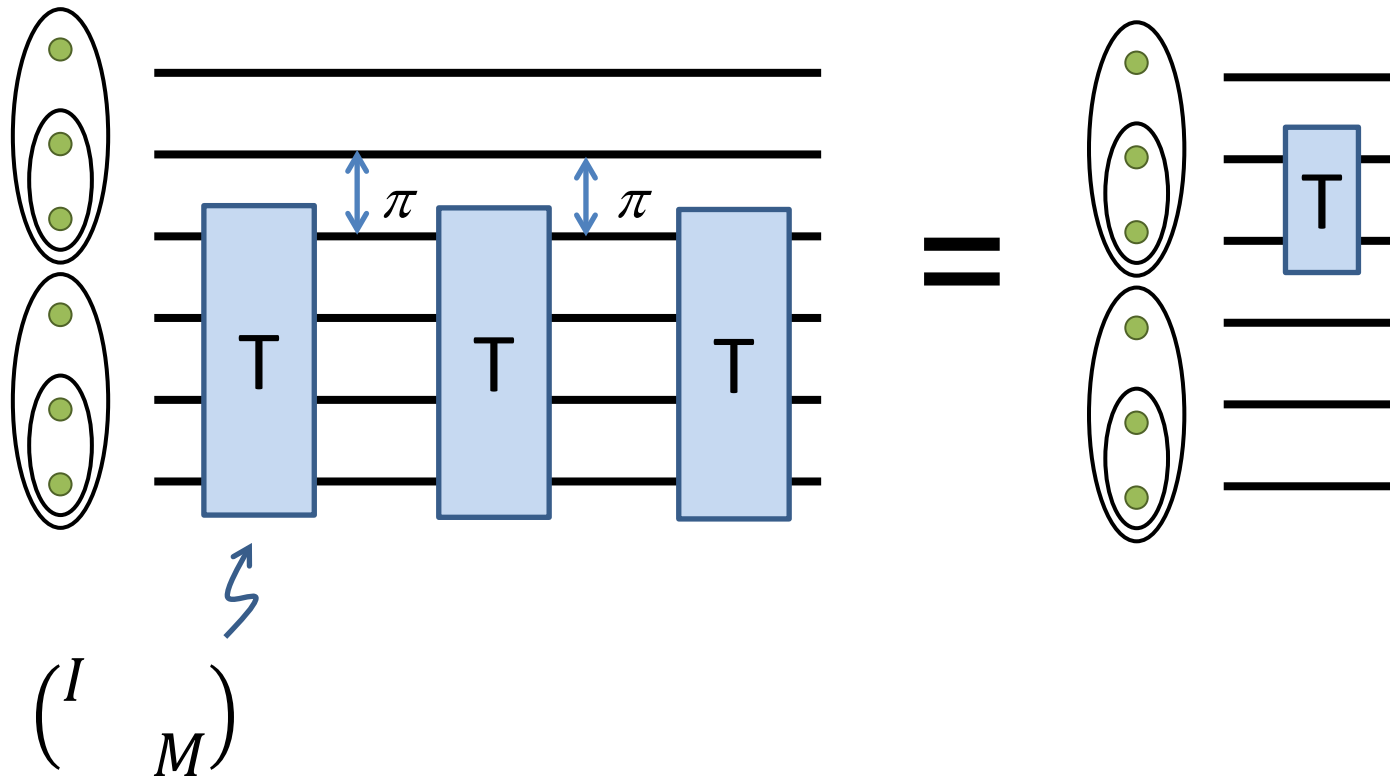
# Sequence Elevation



# Sequence Elevation

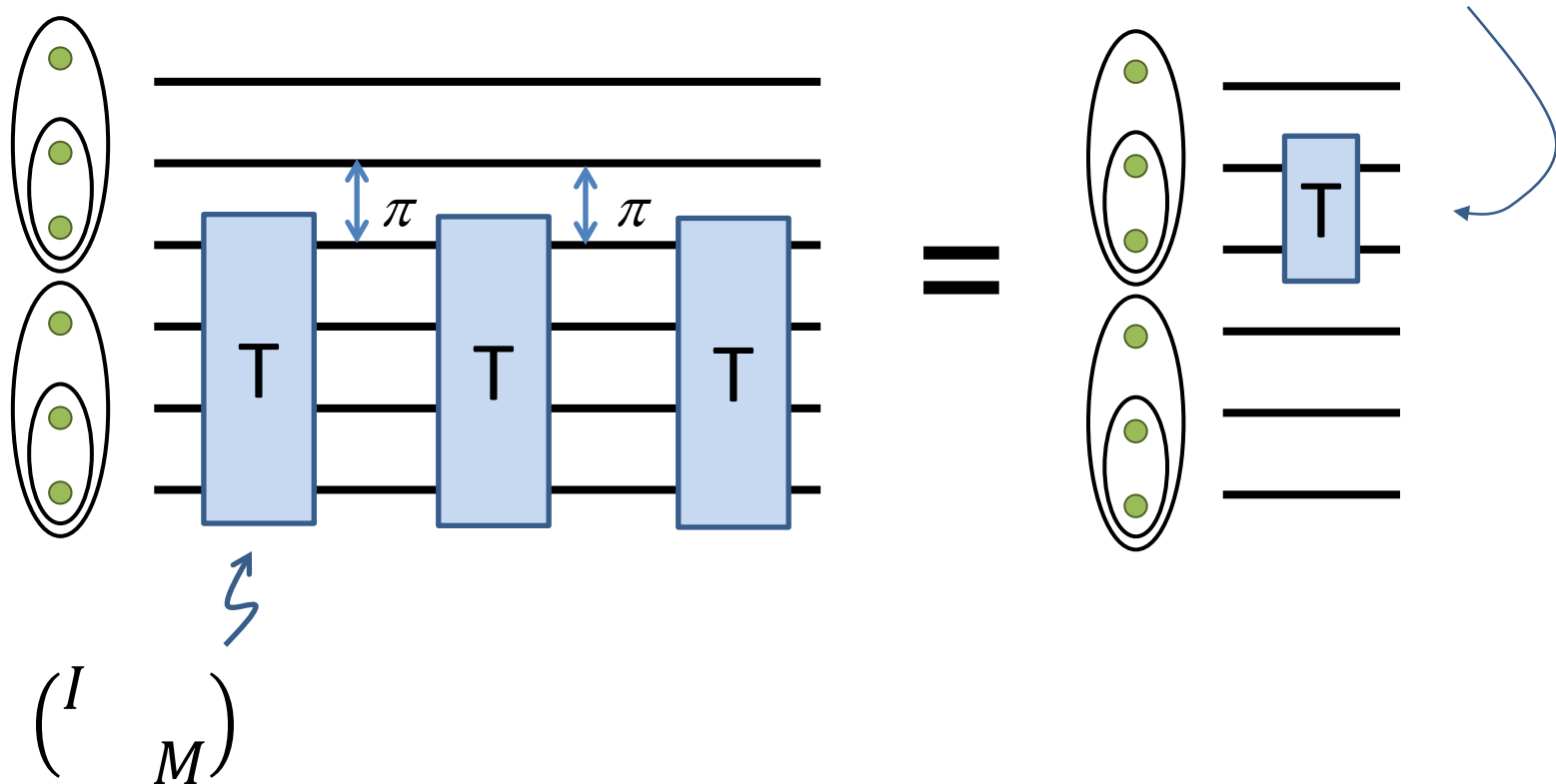


# Sequence Elevation



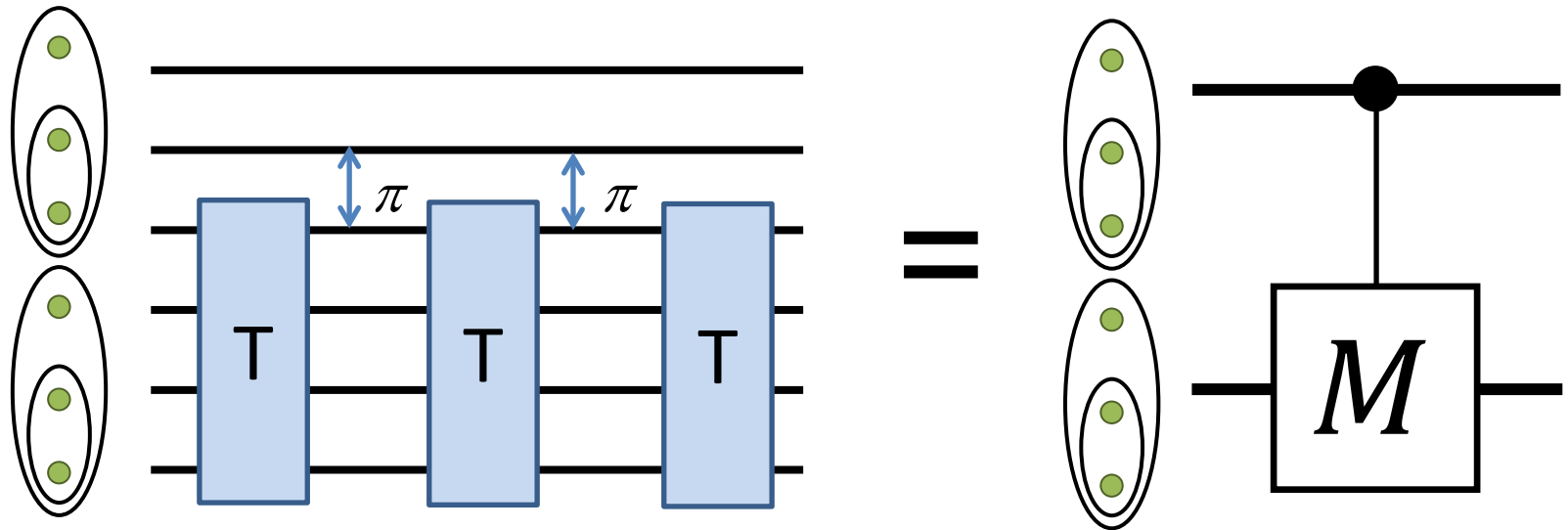
# Sequence Elevation

Applies  $M$  to bottom qubit only if top qubit is in state  $|1\rangle$



# Sequence Elevation

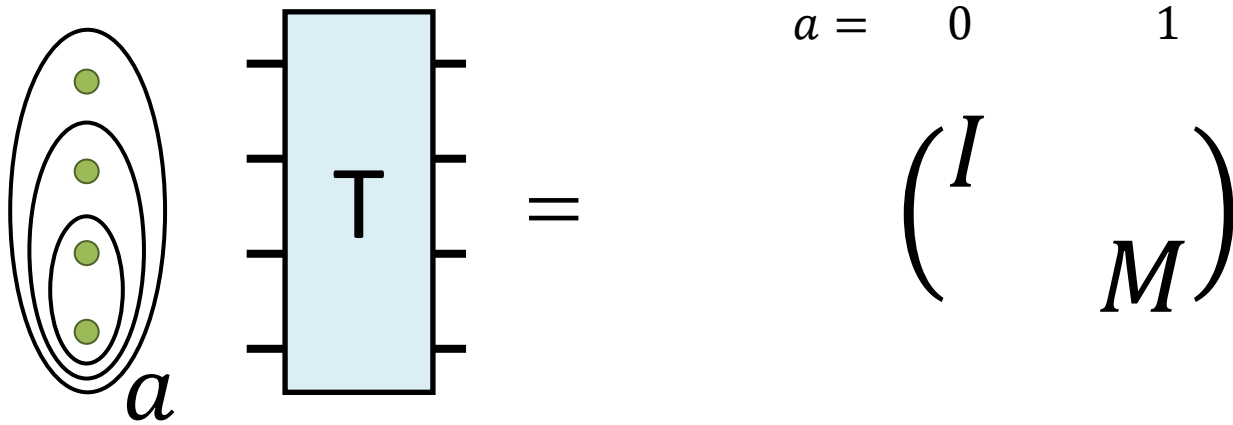
Two-qubit controlled- $M$  gate --- no leakage.



$$M = \hat{\mathbf{n}} \cdot \boldsymbol{\sigma}, \text{ e.g., } M = \underbrace{\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}}_{\text{controlled-NOT}}, \underbrace{\begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}}_{\text{controlled-Phase}}, \dots$$

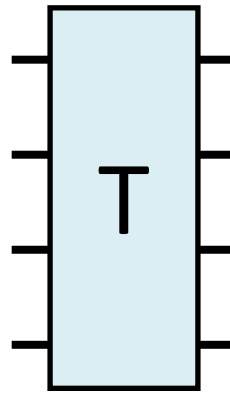
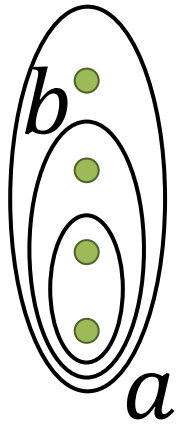
# Constructing T

What we want

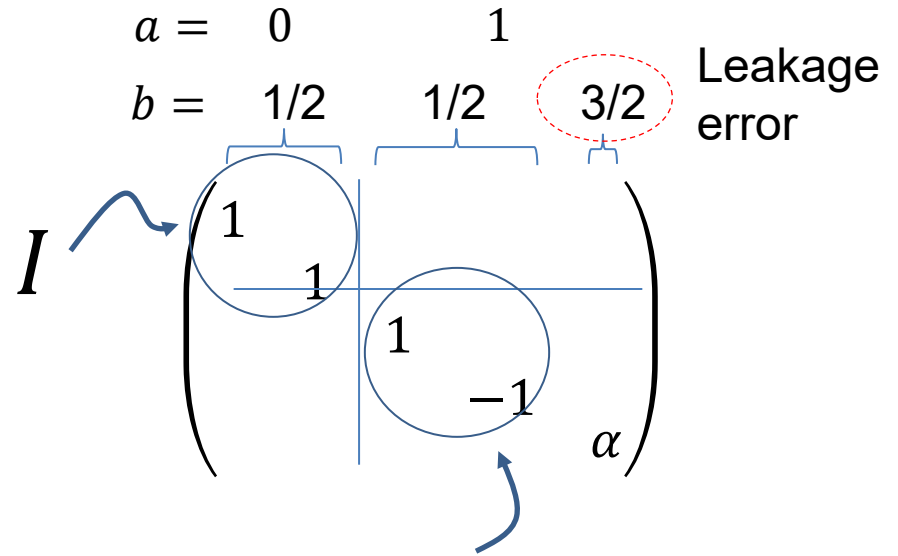


# Constructing T

What we want



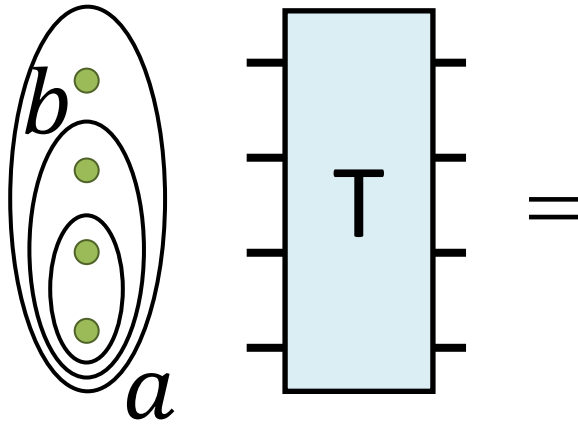
=



$$M = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \text{ for example}$$

# Constructing T

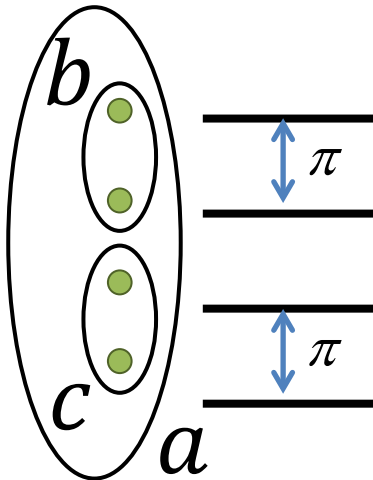
What we want



$$\begin{array}{r}
 a = 0 \qquad \qquad \qquad 1 \\
 c = 1/2 \qquad \qquad 1/2 \qquad \textcircled{3/2}
 \end{array}
 \begin{array}{l}
 \text{Leakage} \\
 \text{error}
 \end{array}$$

$$\left( \begin{array}{c|cc}
 1 & & \\
 \hline
 & 1 & \\
 & & 1 & -1 \\
 & & & \alpha
 \end{array} \right)$$

This is close! (But wrong basis)



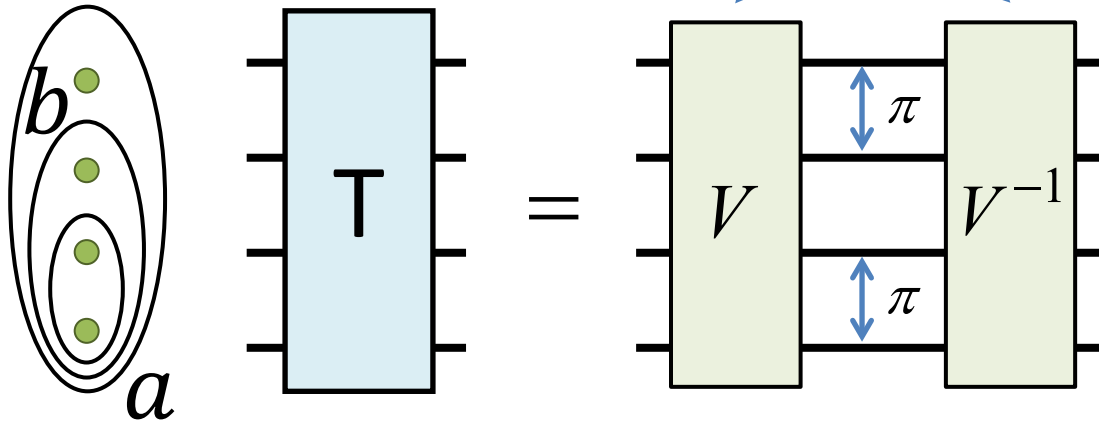
$$\begin{array}{r}
 a = 0 \qquad \qquad \qquad 1 \\
 bc = 00 \ 11 \ 01 \ 10 \ 11
 \end{array}$$

$$\left( \begin{array}{c|ccc}
 1 & & & \\
 \hline
 & 1 & & \\
 & & -1 & \\
 & & & -1 \quad 1
 \end{array} \right)$$

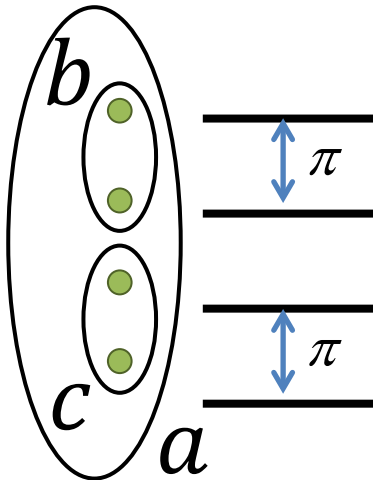


# Constructing T

What we want



This is close! (But wrong basis)

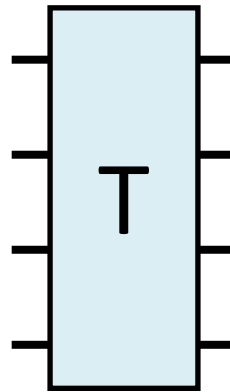
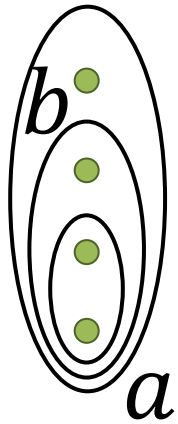


$$a = \begin{array}{cc} 0 & 1 \\ bc = 00 & 11 & 01 & 10 & 11 \end{array}$$

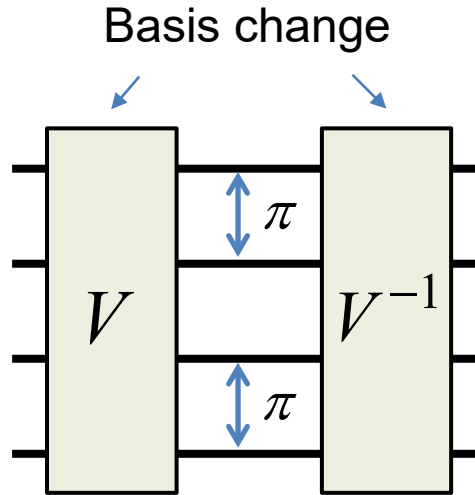
$$\begin{pmatrix} 1 & & & & \\ & 1 & & & \\ \hline & & -1 & & \\ & & & -1 & \\ & & & & 1 \end{pmatrix}$$

# Constructing T

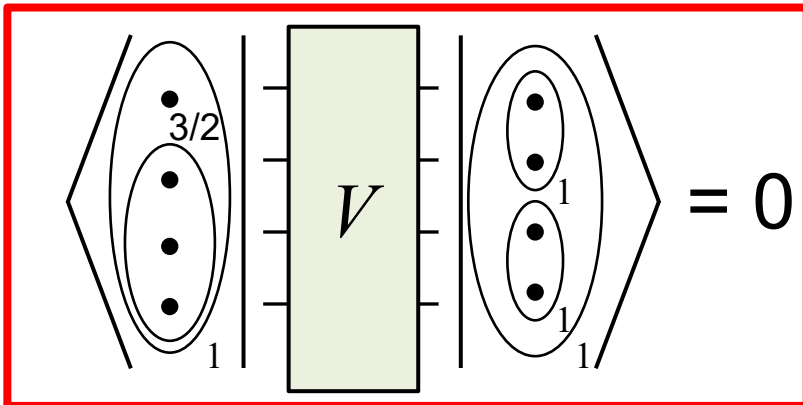
What we want



=

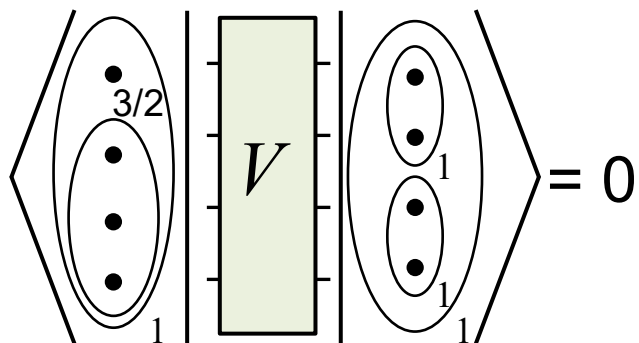


Constraint

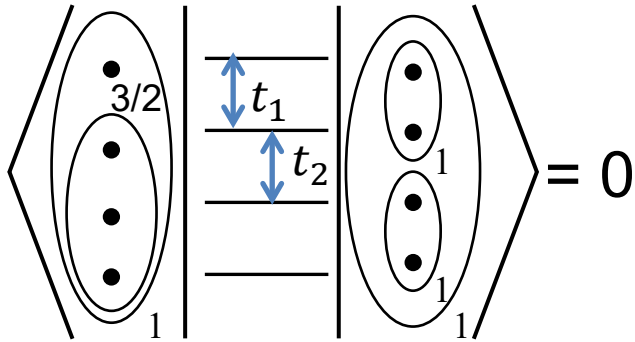


Any  $V$  satisfying this constraint will do the job.

# Satisfying the Constraint



# Satisfying the Constraint



# Satisfying the Constraint

The diagram shows a product of two representations on the left, separated by a vertical line. The first representation is a hexagon containing an inner ellipse with three dots, labeled  $3/2$  and  $1$ . The second representation is a hexagon containing two inner ellipses, each with two dots, labeled  $1$  and  $1$ . Two blue double-headed arrows between the vertical lines are labeled  $t_1$  and  $t_2$ . This is followed by an equals sign and another product of two representations on the right, identical to the ones on the left. To the right of the diagram is the mathematical expression  $\frac{1}{2}(-1 + e^{it_1} + e^{it_2} + e^{i(t_1+t_2)})$ . A blue bracket underlines the entire expression, and a blue arrow points from the bracket down to a red-bordered box.

$$\frac{1}{2}(-1 + e^{it_1} + e^{it_2} + e^{i(t_1+t_2)})$$

Form entirely fixed by  $t_1 = 0, \pi$  and  $t_2 = 0, \pi$  cases

# Satisfying the Constraint

$$\frac{1}{2} (-1 + e^{it_1} + e^{it_2} + e^{i(t_1+t_2)})$$

Form entirely fixed by  $t_1 = 0, \pi$  and  $t_2 = 0, \pi$  cases

For example,  $t_1 = \pi$  and  $t_2 = 0$

$$\frac{1}{2} (-1 + e^{i\pi} + e^{i0} + e^{i(\pi+0)})$$

# Satisfying the Constraint

$$\frac{1}{2} (-1 + e^{it_1} + e^{it_2} + e^{i(t_1+t_2)})$$

Form entirely fixed by  $t_1 = 0, \pi$  and  $t_2 = 0, \pi$  cases

For example,  $t_1 = \pi$  and  $t_2 = 0$

$$\frac{1}{2} (-1 + (-1) + 1 + (-1))$$

# Satisfying the Constraint

$$\frac{1}{2} (-1 + e^{it_1} + e^{it_2} + e^{i(t_1+t_2)})$$

Form entirely fixed by  $t_1 = 0, \pi$  and  $t_2 = 0, \pi$  cases

For example,  $t_1 = \pi$  and  $t_2 = 0$

$$(-1)$$



# Satisfying the Constraint

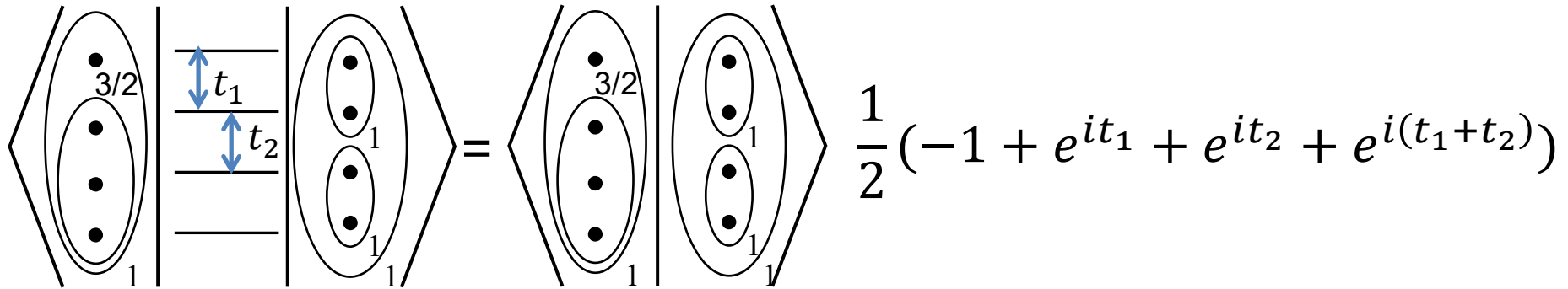
$$\frac{1}{2} (-1 + e^{it_1} + e^{it_2} + e^{i(t_1+t_2)})$$

Form entirely fixed by  $t_1 = 0, \pi$  and  $t_2 = 0, \pi$  cases

For example,  $t_1 = \pi$  and  $t_2 = 0$

$$(-1)$$

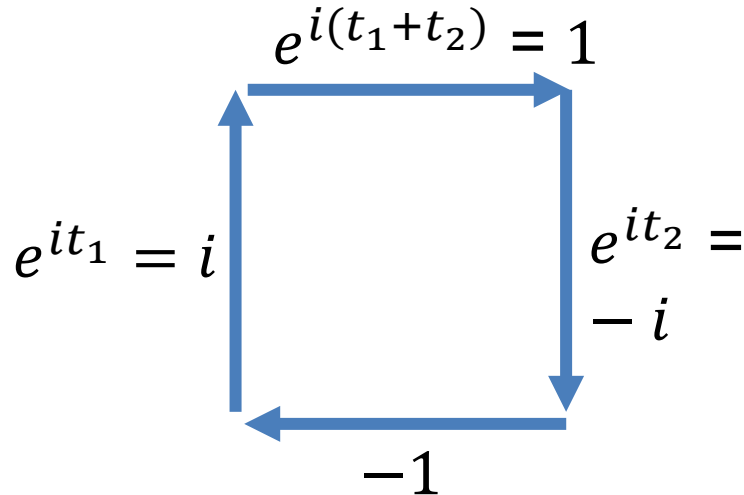
# Satisfying the Constraint



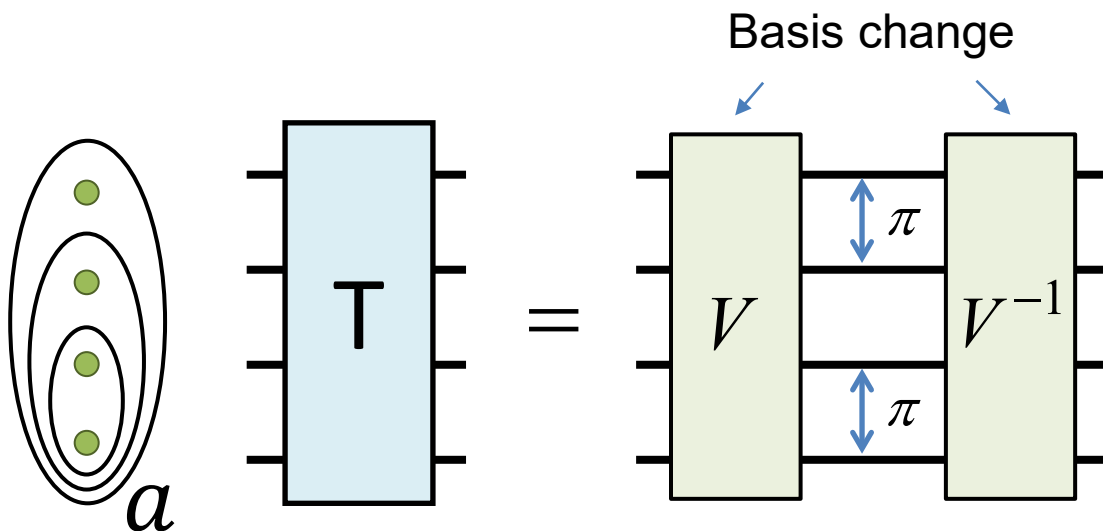
Need destructive interference

$$t_1 = \frac{\pi}{2}$$

$$t_2 = \frac{3\pi}{2}$$

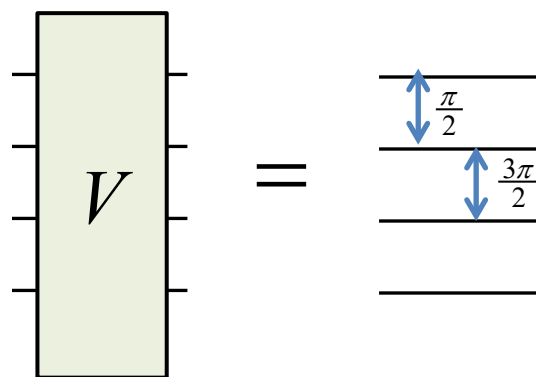
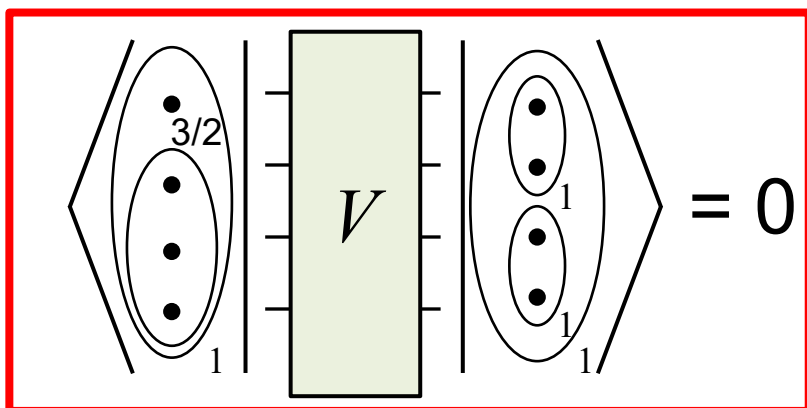


# Constructing T

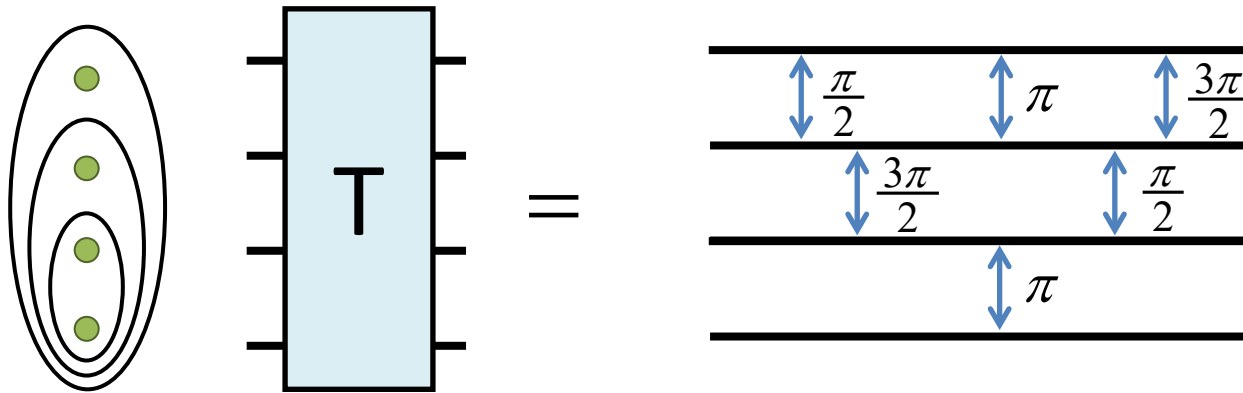


Constraint

Optimal Solution

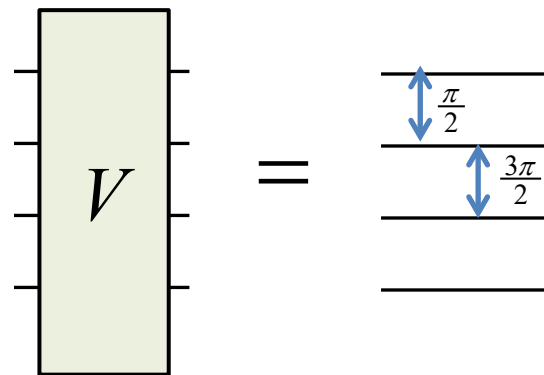
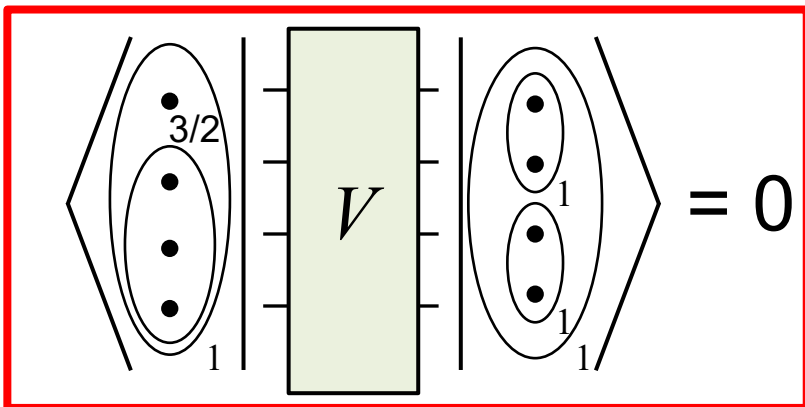


# Constructing T

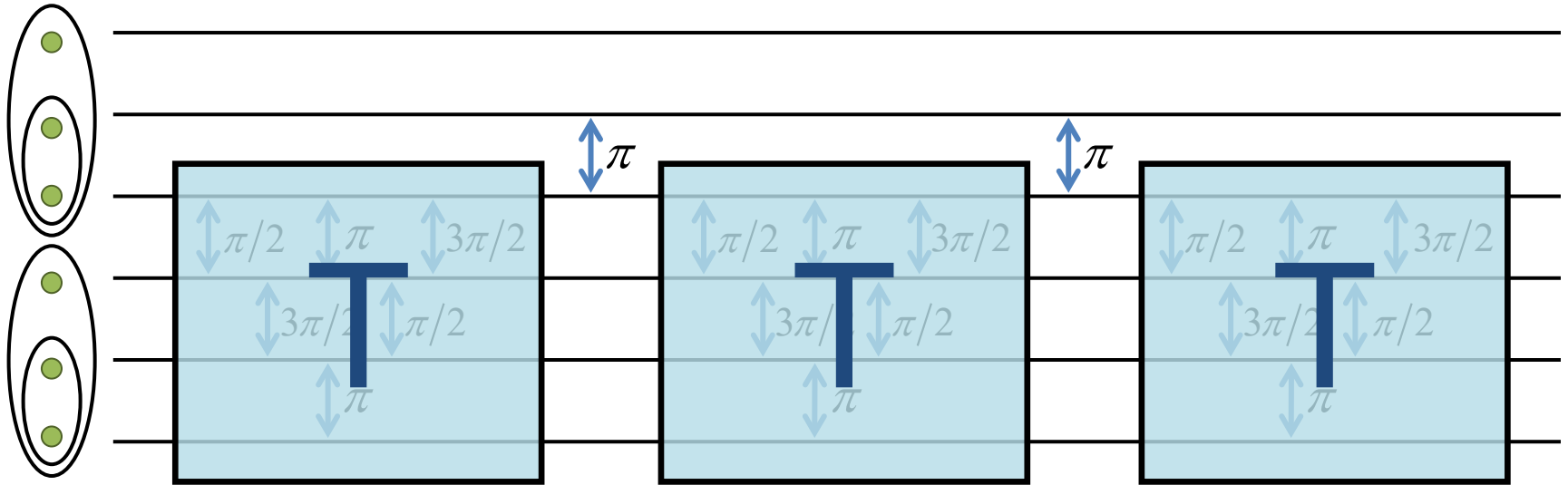


Constraint

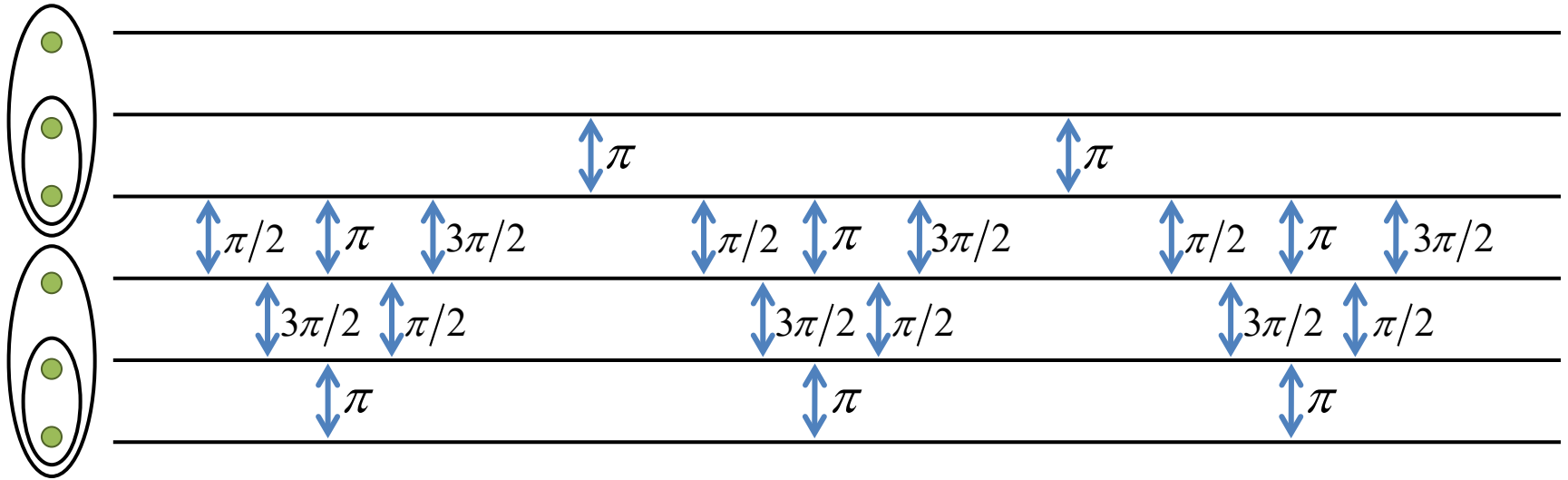
Optimal Solution



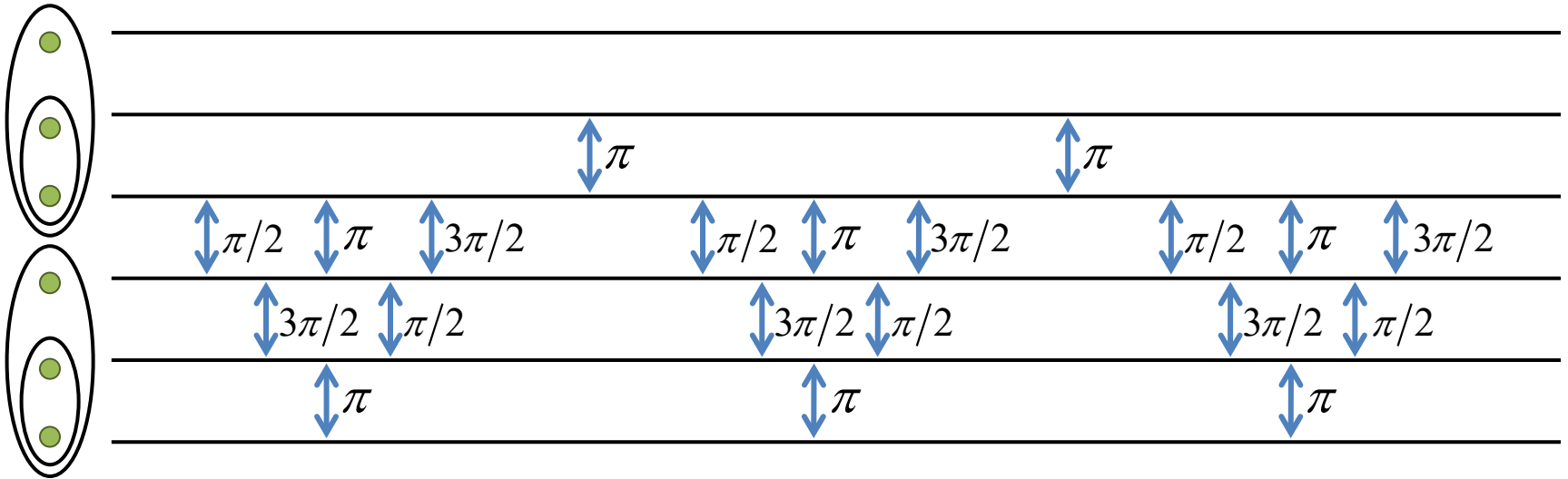
# Full Sequence



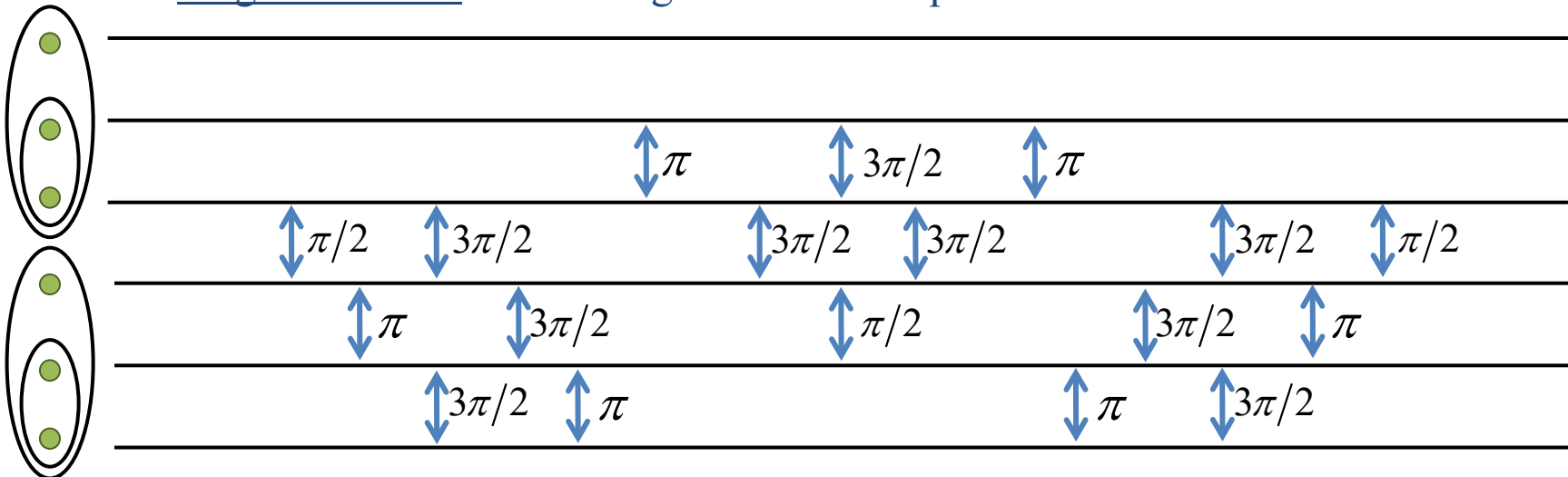
# Full Sequence



# Full Sequence



- Original version of the Fong-Wandzura Sequence



PHYSICAL REVIEW A **93**, 010303(R) (2016)

## Simple derivation of the Fong-Wandzura pulse sequence

Daniel Zeuch and N. E. Bonesteel

*Department of Physics and National High Magnetic Field Laboratory, Florida State University, Tallahassee, Florida 32310, USA*

(Received 4 September 2015; published 25 January 2016)

We give an analytic construction of a class of two-qubit gate pulse sequences that act on five of the six spin- $\frac{1}{2}$  particles used to encode a pair of exchange-only three-spin qubits. Within this class, the problem of gate construction reduces to that of finding a smaller sequence that acts on four spins and is subject to a simple constraint. The optimal sequence satisfying this constraint yields a two-qubit gate sequence equivalent to that found numerically by Fong and Wandzura. Our construction is sufficiently simple that it can be carried out entirely with pen, paper, and knowledge of a few basic facts about quantum spin.

DOI: [10.1103/PhysRevA.93.010303](https://doi.org/10.1103/PhysRevA.93.010303)



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## Efficient two-qubit pulse sequences beyond CNOT

D. Zeuch <sup>1</sup> and N. E. Bonesteel <sup>2</sup>

<sup>1</sup>*Peter Grünberg Institute, Theoretical Nanoelectronics, Forschungszentrum Jülich, D-52425 Jülich, Germany*

<sup>2</sup>*Department of Physics and National High Magnetic Field Laboratory, Florida State University, Tallahassee, Florida 32310, USA*



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We design efficient controlled-rotation gates with arbitrary angle acting on three-spin encoded qubits for exchange-only quantum computation. Two pulse sequence constructions are given. The first is motivated by an analytic derivation of the efficient Fong-Wandzura sequence for an exact CNOT gate. This derivation, briefly reviewed here, is based on elevating short sequences of SWAP pulses to an entangling two-qubit gate. To go beyond CNOT, we apply a similar elevation to a modified short sequence consisting of SWAPS and one pulse of arbitrary duration. This results in two-qubit sequences that carry out controlled-rotation gates of arbitrary angle. The second construction streamlines a class of arbitrary CPHASE gates established earlier. Both constructions are based on building two-qubit sequences out of subsequences with special properties that render each step of the construction analytically tractable.

DOI: [10.1103/PhysRevB.102.075311](https://doi.org/10.1103/PhysRevB.102.075311)

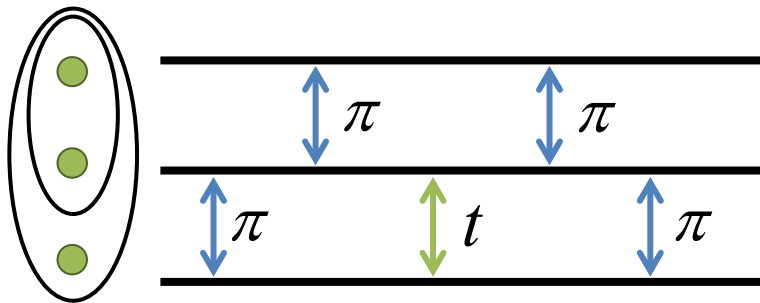
# New Sequences?

---

Strategy: “Elevate” simple three-spin pulse sequence to a two-qubit gate.

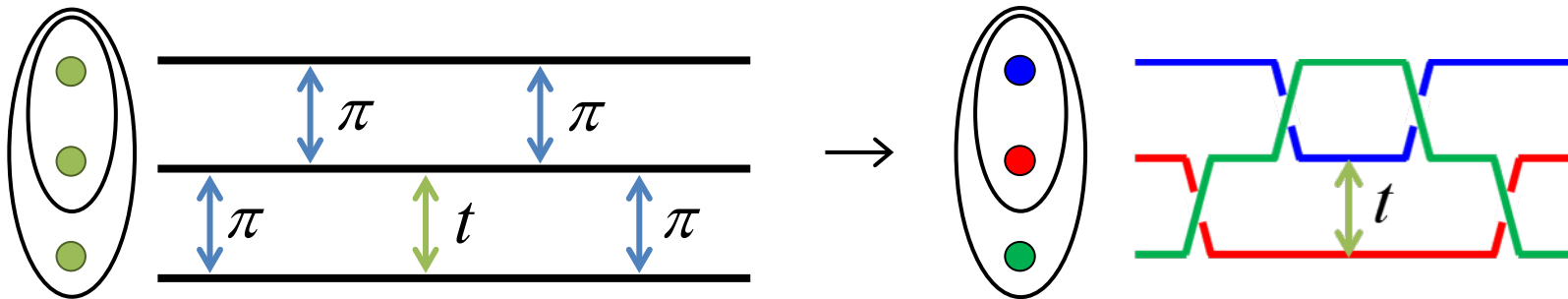
# Another Simple Sequence

Strategy: “Elevate” simple three-spin pulse sequence to a two-qubit gate.



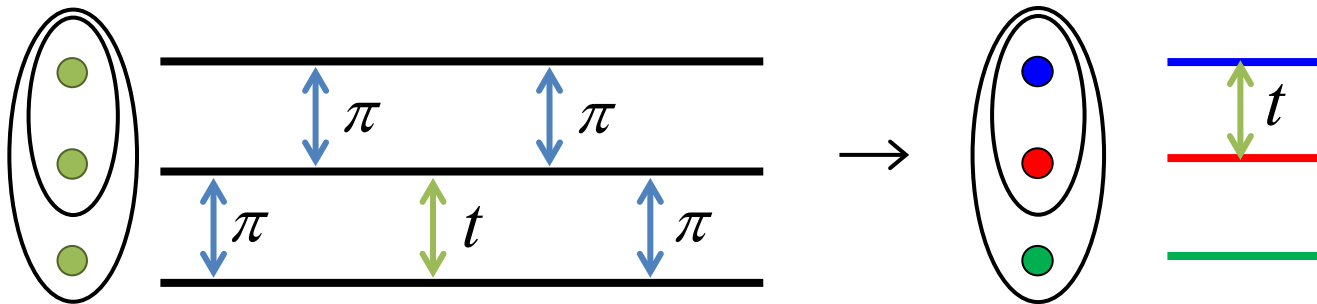
# Another Simple Sequence

Strategy: “Elevate” simple three-spin pulse sequence to a two-qubit gate.



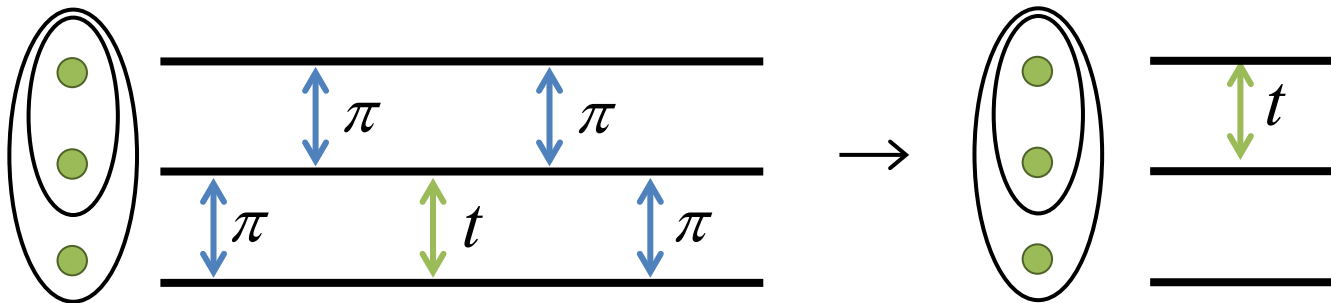
# Another Simple Sequence

Strategy: “Elevate” simple three-spin pulse sequence to a two-qubit gate.



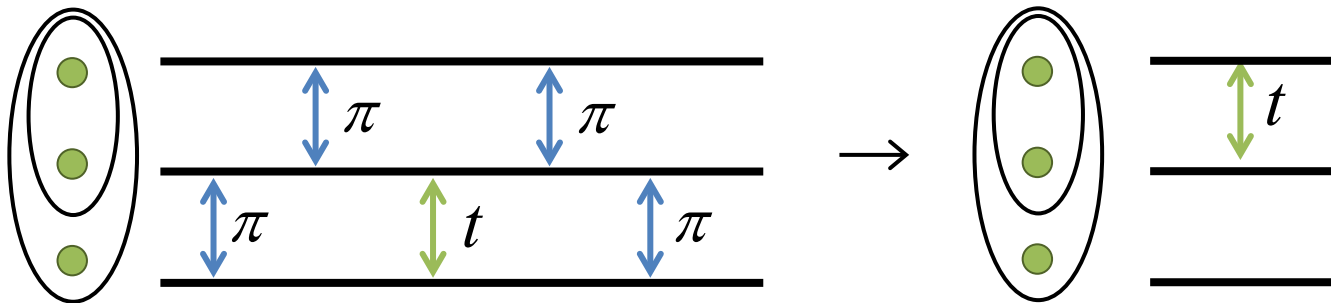
# Another Simple Sequence

Strategy: “Elevate” simple three-spin pulse sequence to a two-qubit gate.



# Sequence Identity

Strategy: “Elevate” simple three-spin pulse sequence to a two-qubit gate.

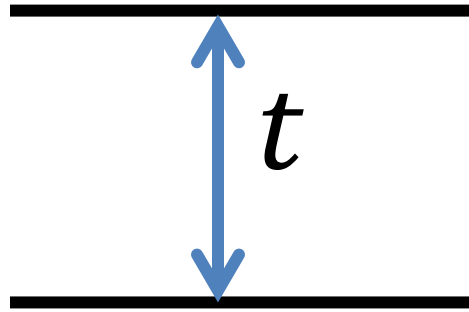
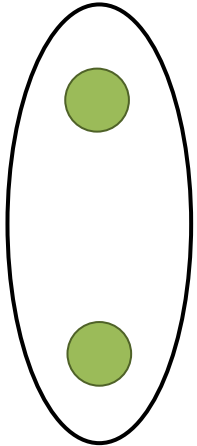


$$\begin{pmatrix} 1 & \\ & e^{-it} \end{pmatrix}$$

Identity holds for *any*  $t$ , not just  $t=0, \pi$

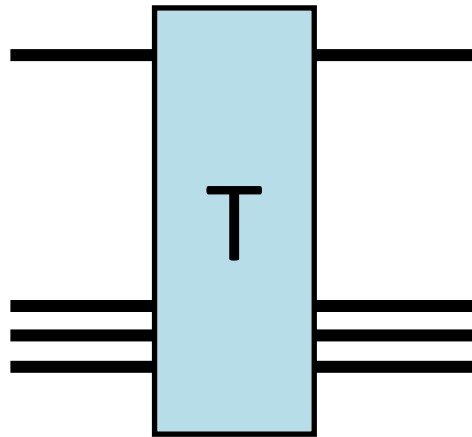
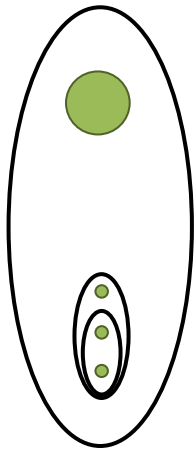


# “Elevating” $t$ Pulse



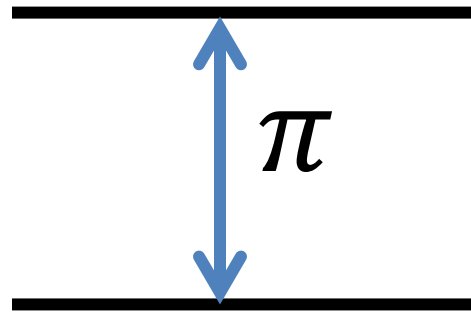
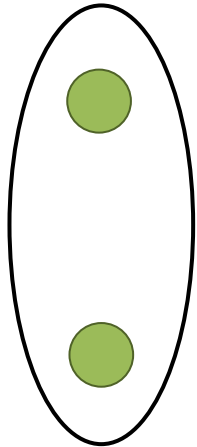
$$\begin{pmatrix} 1 & \\ & e^{it} \end{pmatrix}$$

$M$  can be any 2x2 unitary.  
No longer require  $M^2 = I$

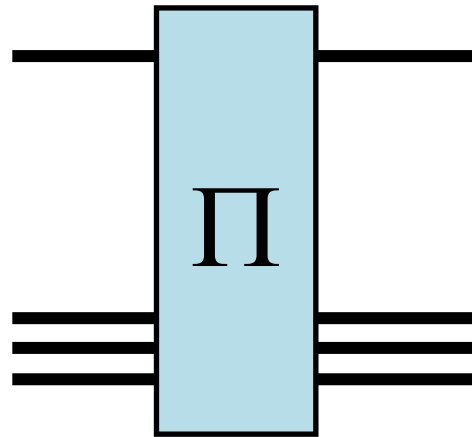
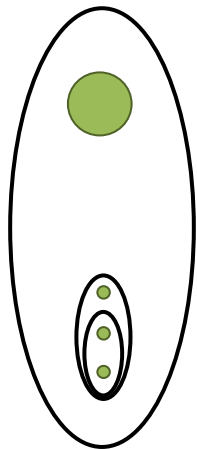


$$\begin{pmatrix} I & \\ & M \end{pmatrix}$$

# “Elevating” *SWAP* Pulse

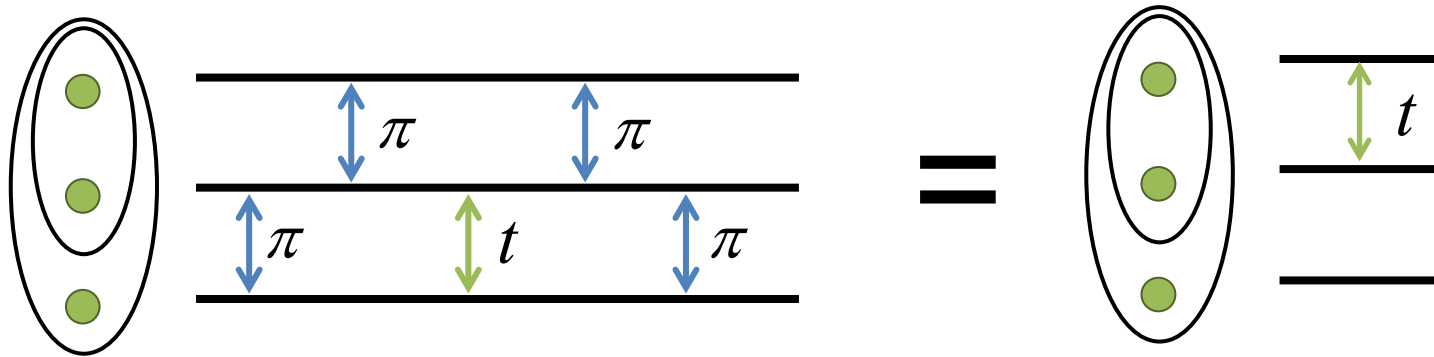


$$\begin{pmatrix} 1 & \\ & -1 \end{pmatrix}$$

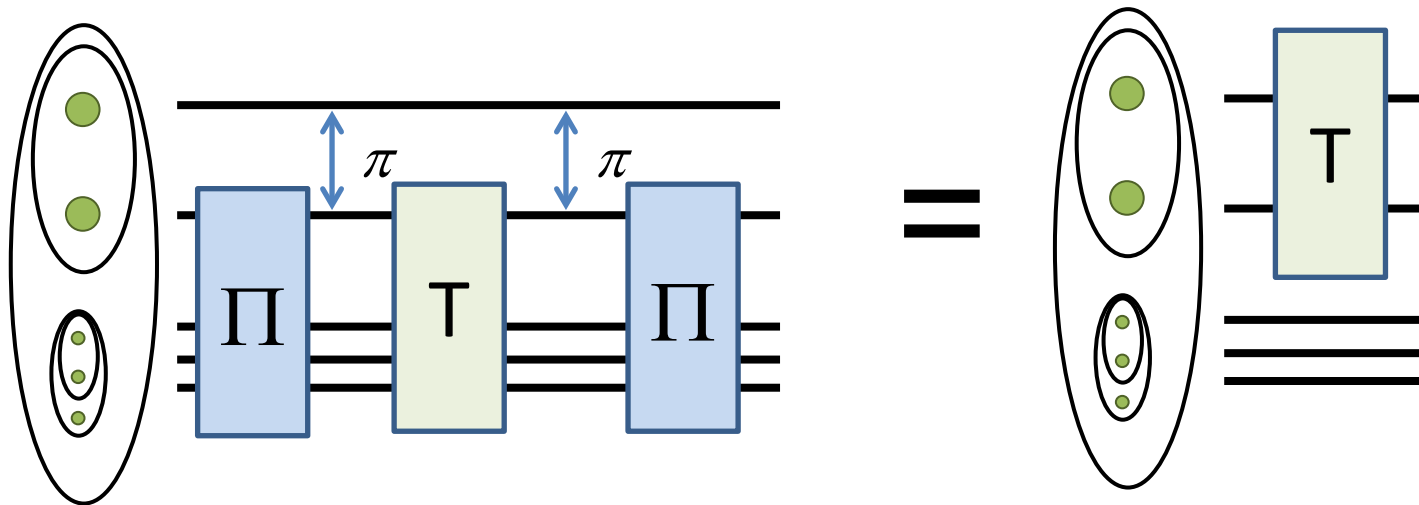


$$\begin{pmatrix} I & \\ & -I \end{pmatrix}$$

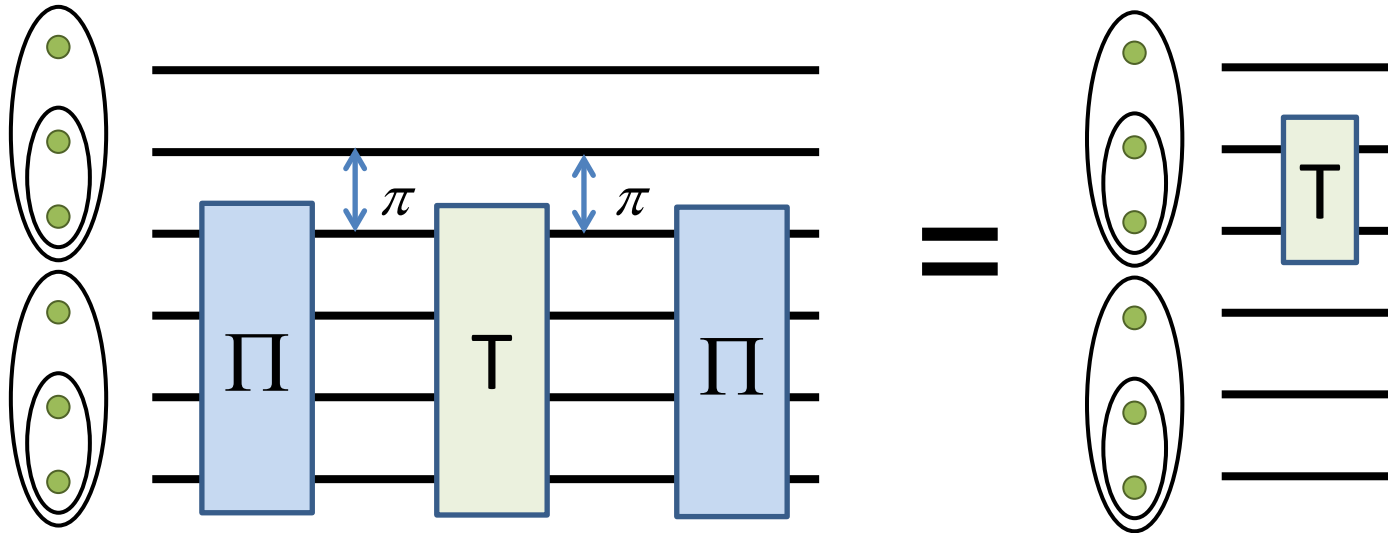
# Sequence Elevation



# Sequence Elevation

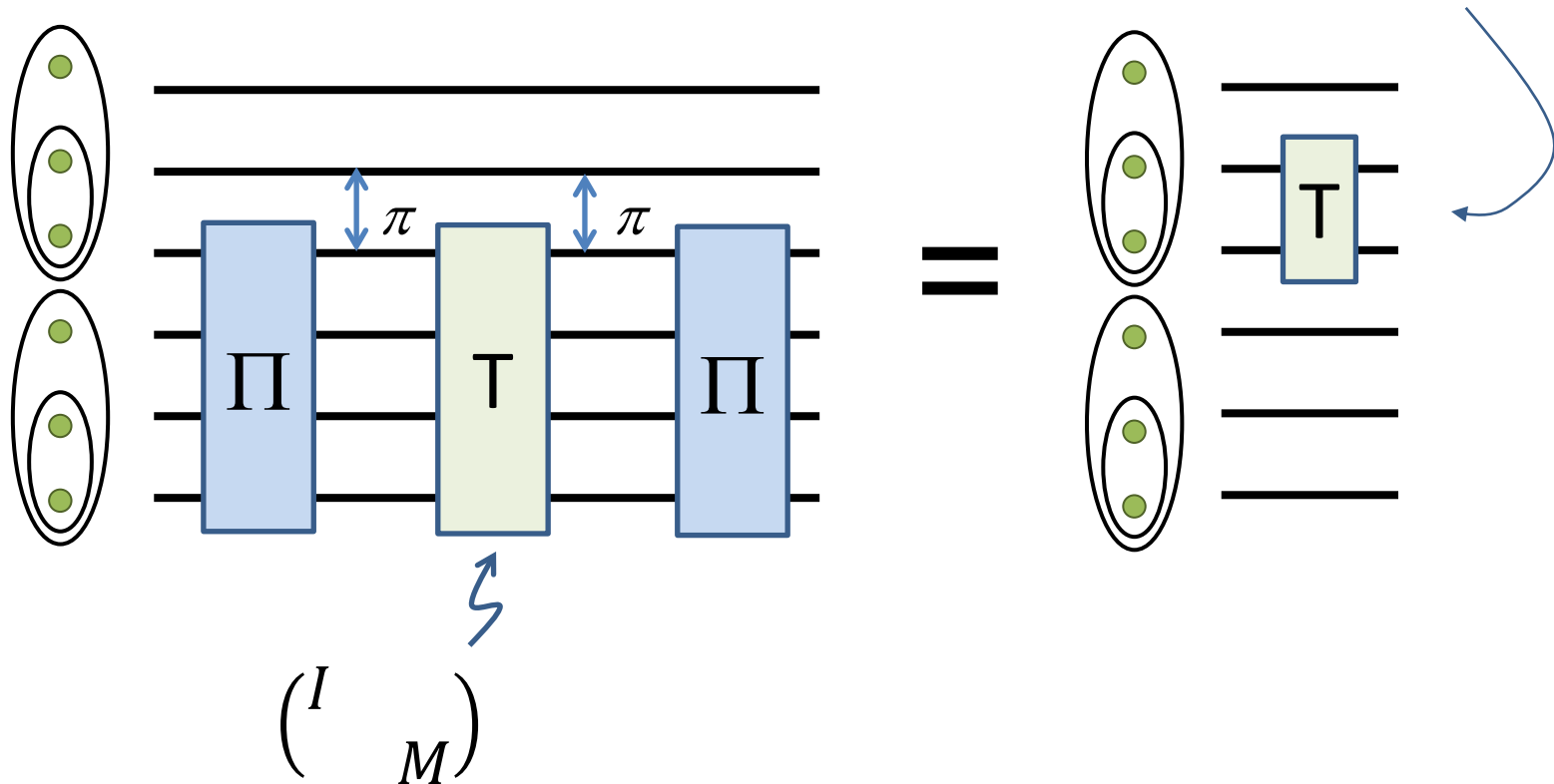


# Sequence Elevation

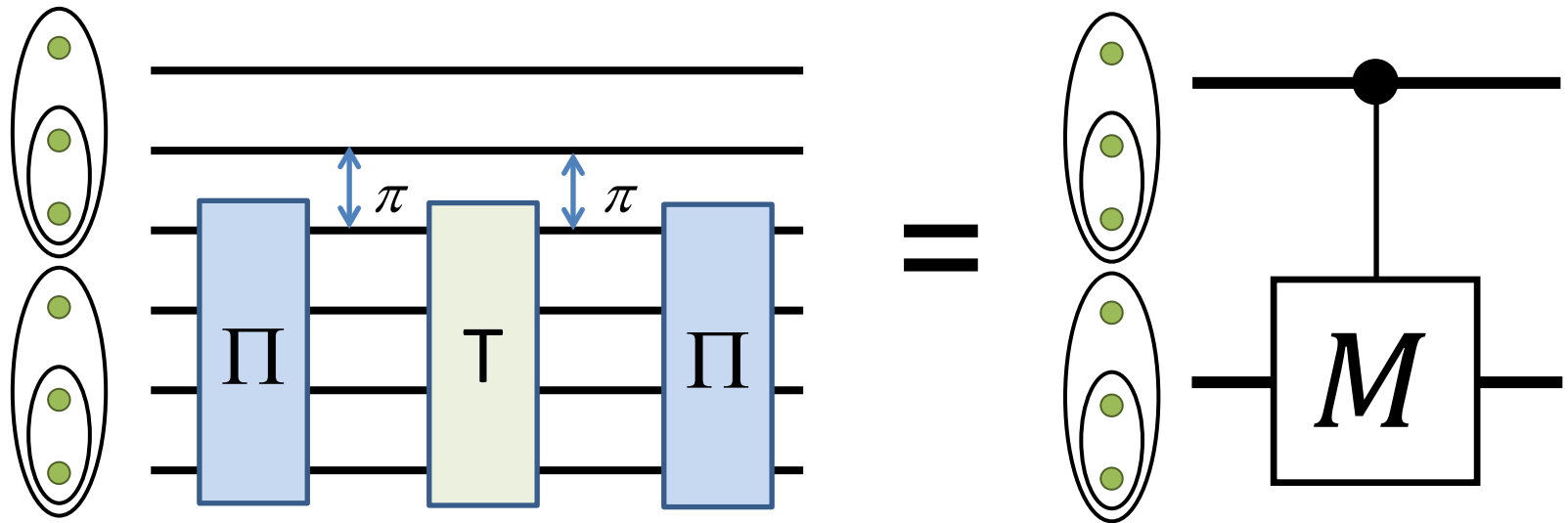


# Sequence Elevation

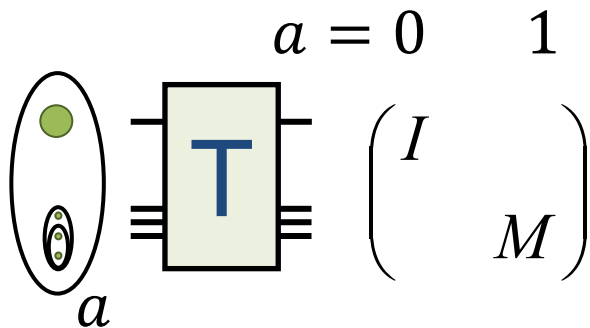
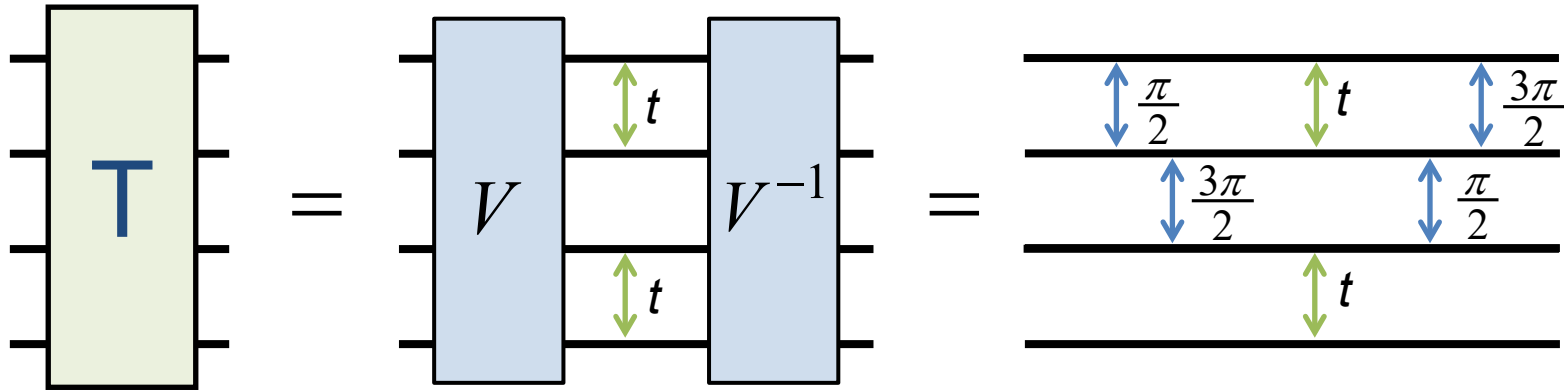
Applies  $M$  to bottom qubit only if top qubit is in state  $|1\rangle$



# Sequence Elevation



# Constructing T



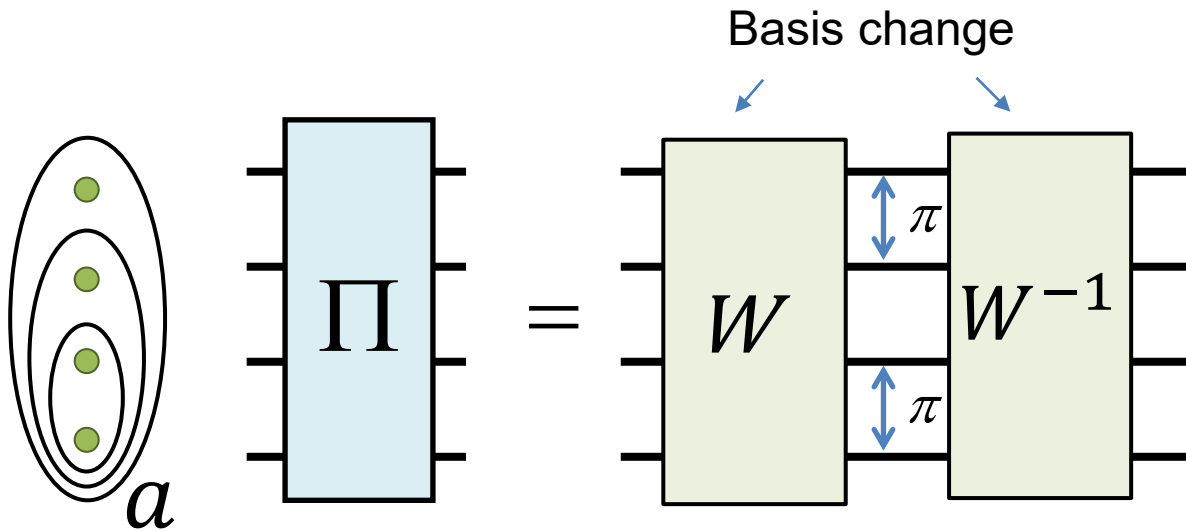
$$M(t) = e^{i \xi(t)} e^{i \phi(t) \hat{n}(t) \cdot \sigma / 2}$$

$$\phi(t) = 2 \arccos\left(\frac{5 \cos(t/2) + 3 \cos(3t/2)}{8}\right)$$

$$\xi(t) = -t/2$$

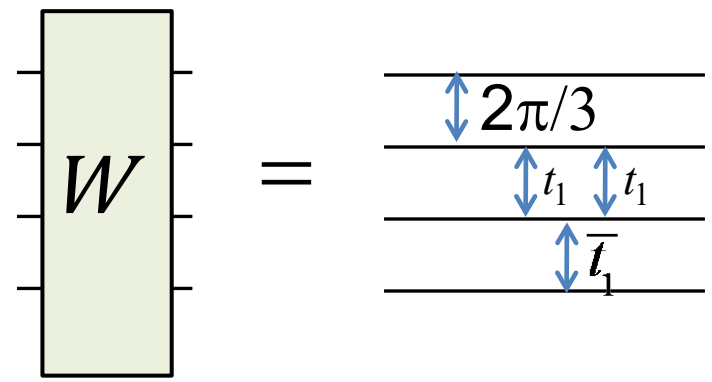
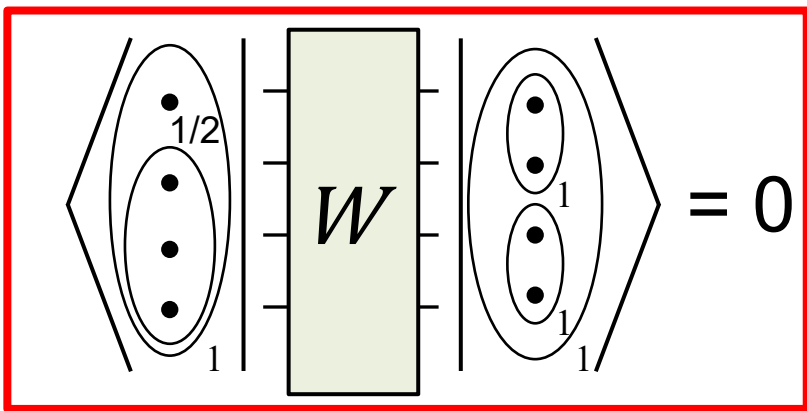


# Constructing $\Pi$

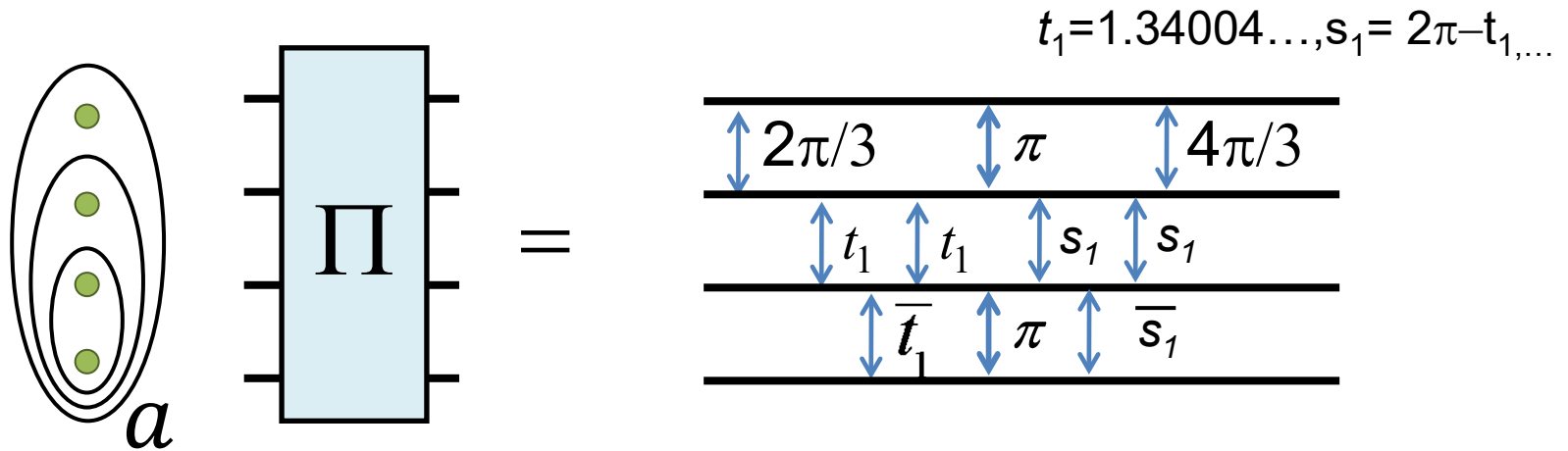


Constraint

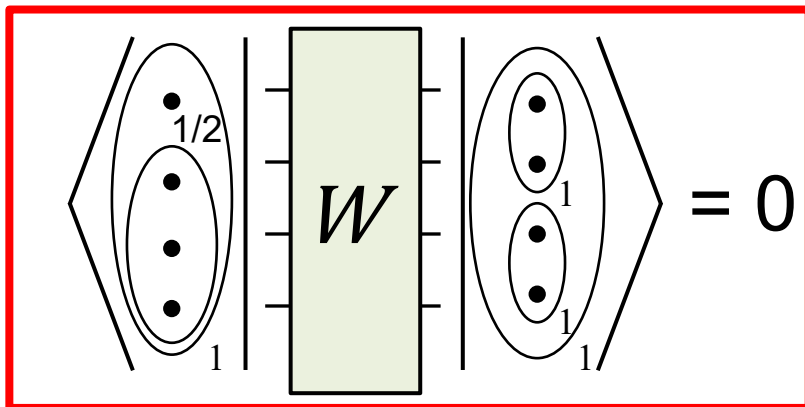
Solution



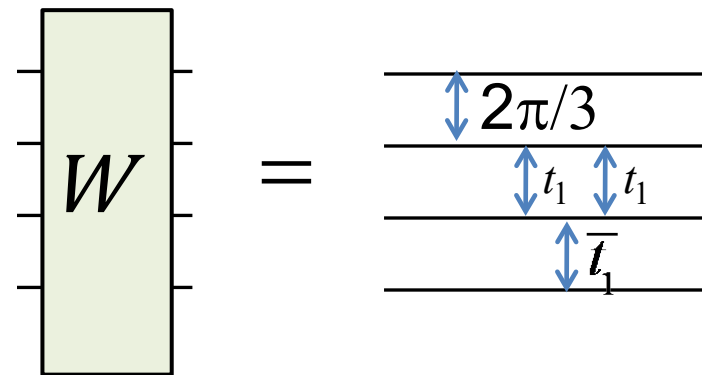
# Constructing $\Pi$



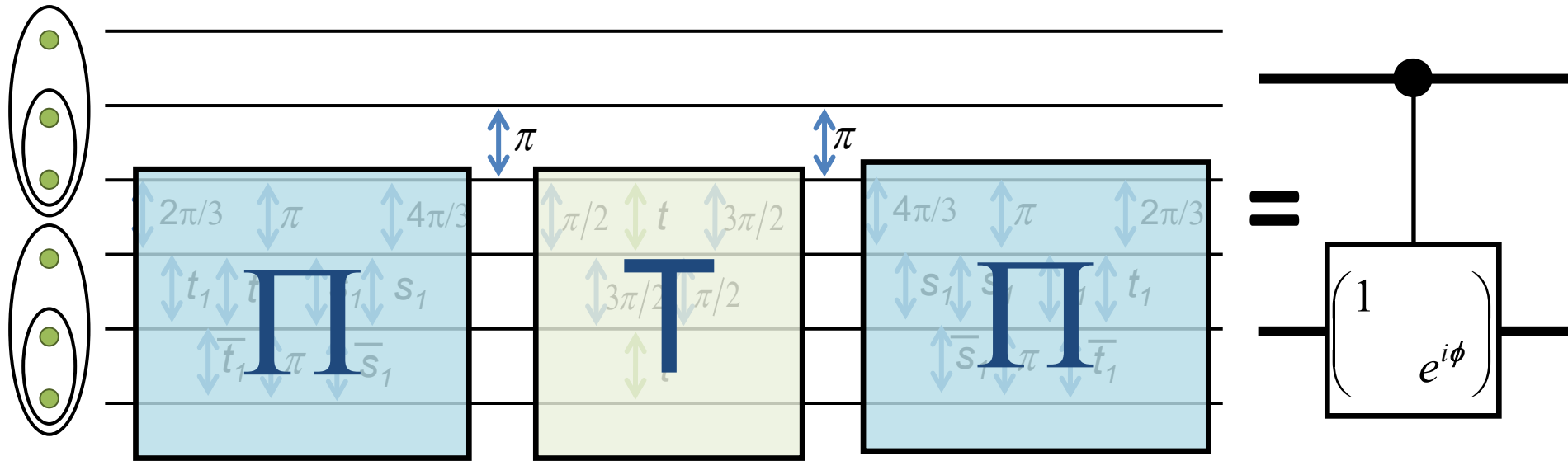
Constraint



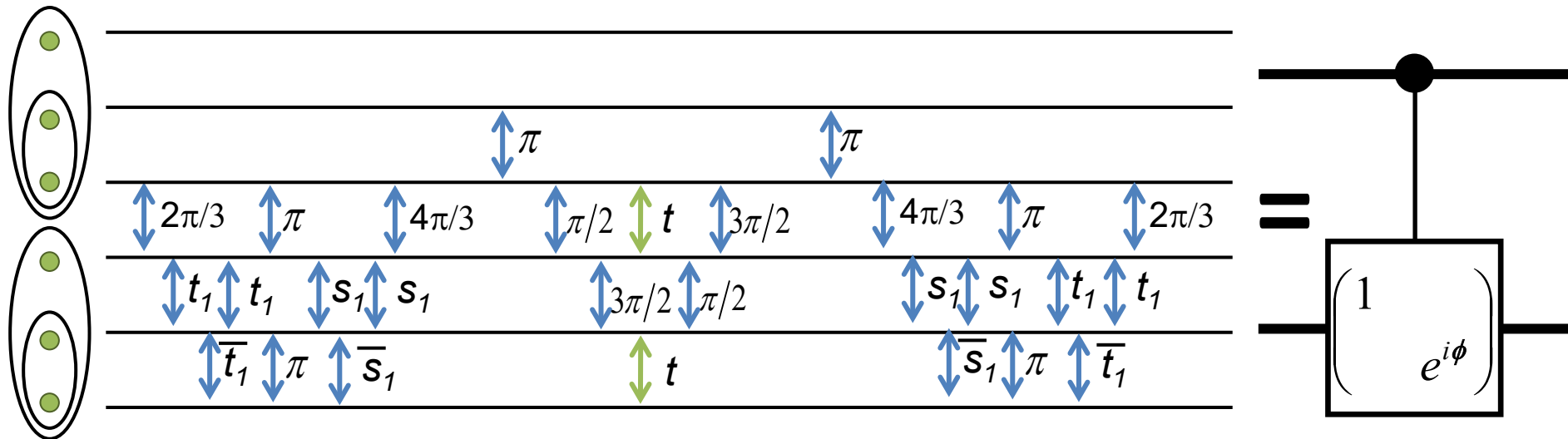
Solution



# Full Sequence



# Full Sequence

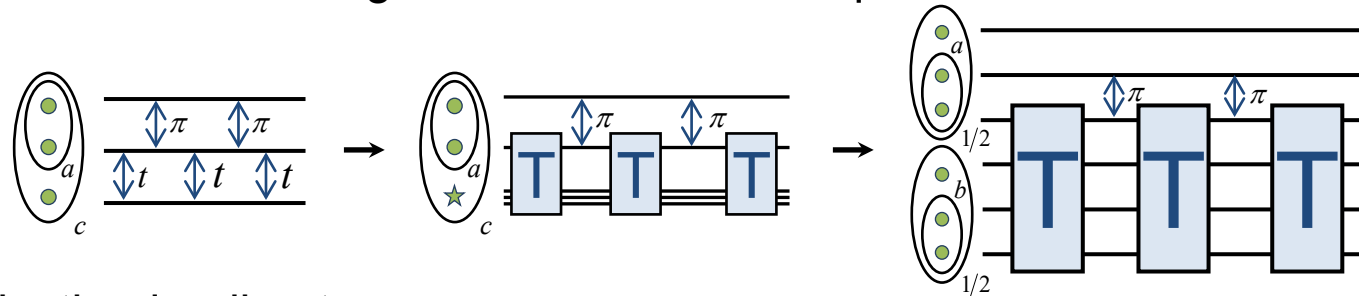


$$t_1 = 1.34004\dots, s_1 = 2\pi - t_1, \dots$$

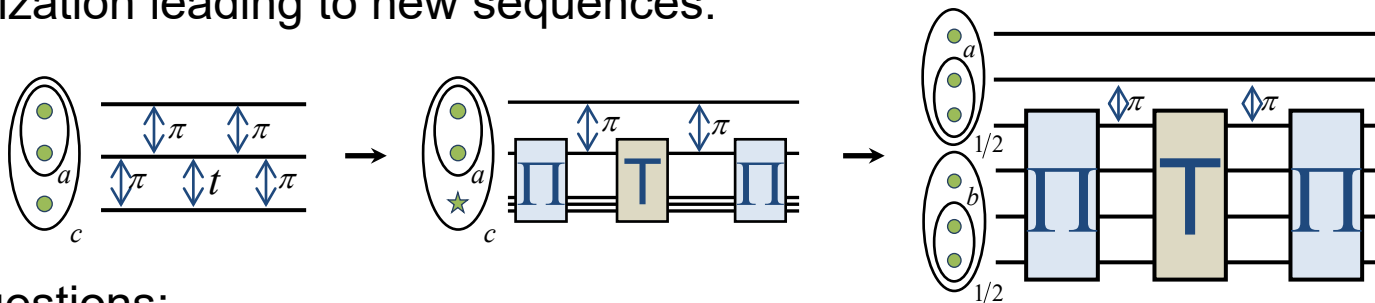
$$\phi(t) = 2 \arccos\left(\frac{5 \cos(t/2) + 3 \cos(3t/2)}{8}\right)$$

# Summary

Analytic Derivation of Fong-Wandzura CNOT sequence:



Generalization leading to new sequences:



Open questions:

- 1) Can we prove Fong-Wandzura sequence is truly optimal?
- 2) More efficient general gate constructions?
- 3) Can these tools be used to construct more “robust” sequences?

Zeuch, NEB, Phys. Rev. A **98**, 010303 (2016)

Zeuch, NEB, Phys. Rev. B **102**, 075311 (2020)